

NVRAM

Peripherals No More

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Oldsmobiles

- GM
- Manufactured 1897~2004
- Went down hill in 1980's
- Desperation to survive



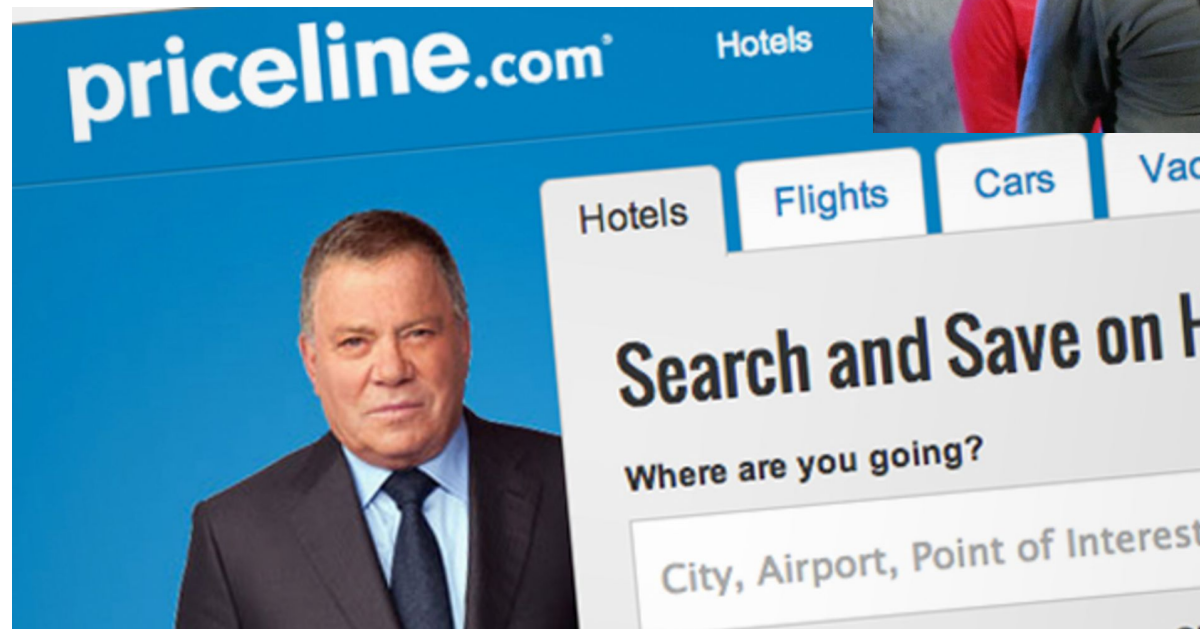
A portrait of actress Melanie Shatner, smiling, with her hands on her hips. She is wearing a dark, short-sleeved top. The background is dark and out of focus.

MELANIE SHATNER



William Shatner

NVRAM





NOT your father's Oldsmobile

- Fond memories
- New Generation of Oldsmobile
- Not your father's Oldsmobile



1st FISS Workshop

- International Workshop on File and Storage Systems
- Organized by KIISE SIGFAST



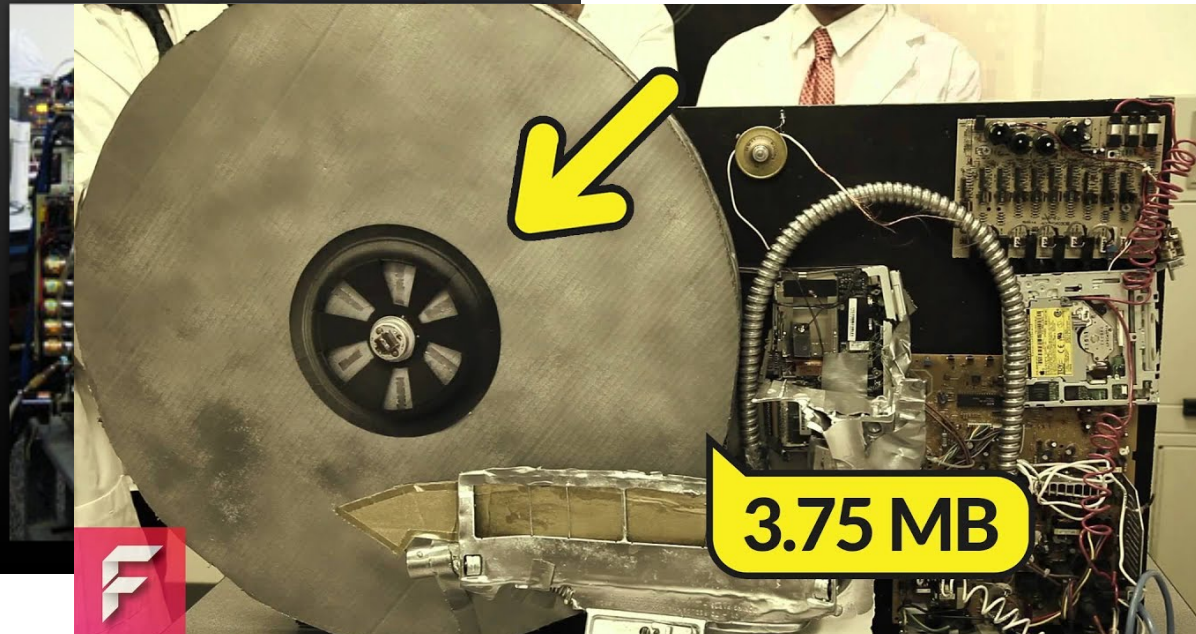
1st FISS Workshop

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Hard Disk



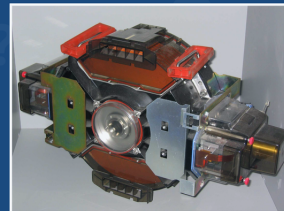
YES! your father's Oldsmobile Storage Device



YES! your father's Oldsmobile Storage Device

did you know ?

the first 1GB hard disk drive was announced in 1980 which weighed about 550 pounds, and had a price tag of \$40,000?



Floppy disks

- distribute software
- transfer files
- back-up data



Disk data storage milestones

- 1971: first 8" floppy disk, IBM



- 1991: first 1GB hard disk drive, IBM



- 2000: 1 inch disk drive, IBM

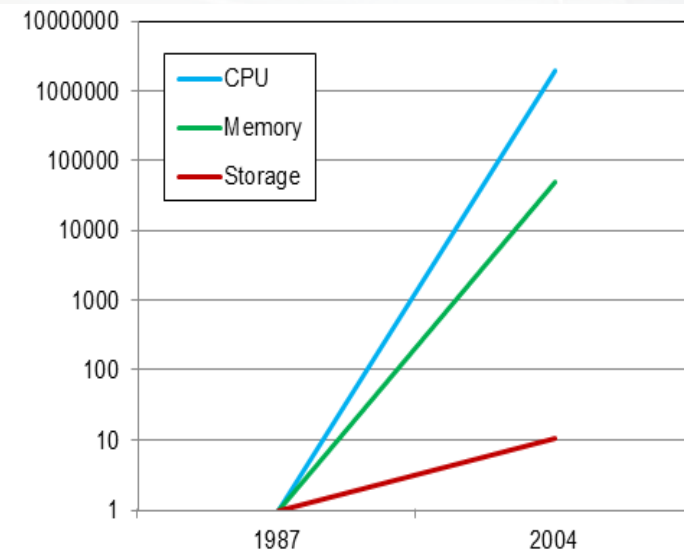
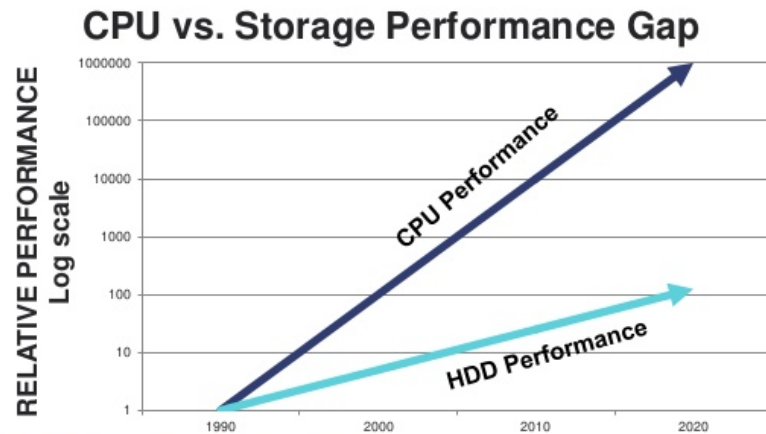


UNIVERSITÉ
PAUL-ÉLIE DUBREUIL

Jean-Pierre Binger-Gaspard - 2010-11-10-2009

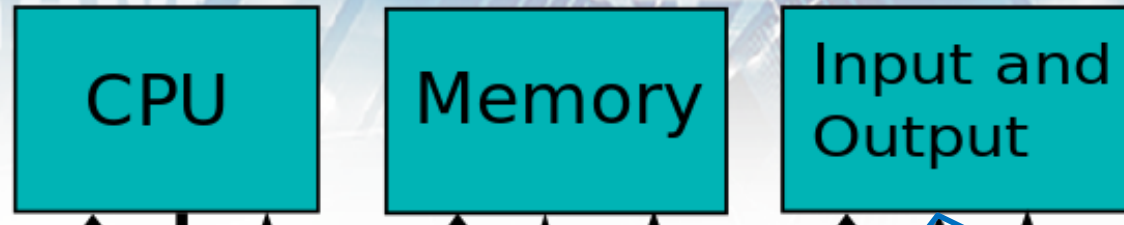
UFR EILA
ÉTUDES INTERCULTURELLES
DE LANGUES APPLIQUÉES

YES! your father's Oldsmobile Storage Device



	1987	2004	Increase Multiple
CPU Performance	1 MIPS	2,000,000 MIPS	2,000,000 x
Memory Performance	100 usec	2 nsec	50,000 x
Disk Drive Performance	60 msec	5.3 msec	11x

YES! your father's ~~Oldsmobile~~ Storage Device



Types of Peripheral Devices

Input



Output



Storage



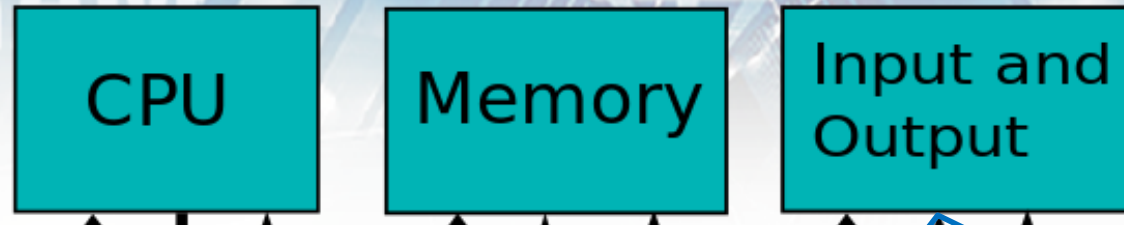
Communication



Rest of the gang!

Peripheral: Auxiliary, Supplementary, relating to periphery

YES! your father's ~~Oldsmobile~~ Storage Device



Types of Peripheral Devices



Peripheral: Auxiliary, Supplementary, relating to periphery

NOT your father's Oldsmobile Storage Device

- **New generation of storage**
 - Ultra Low Latency (ULL) drives
 - NVMe



	Samsung Z-SSD (SZ985)	Intel Optane (P4800X)
Technology	Z-NAND	3D Xpoint
Capacity	800GB	750GB
Sequential Read/Write (GB/s)	3.2GB/s (Both)	2.4GB/s Read 2GB/s Write
Random Read/Write (IOPS)	750K Read 170K Write	550K Read 500K Write
Random Read Latency	12-20us	10us
Random Write Latency	16us	10us



NOT your father's Oldsmobile Storage Device

- New generation of storage
 - DIMM slotted storage

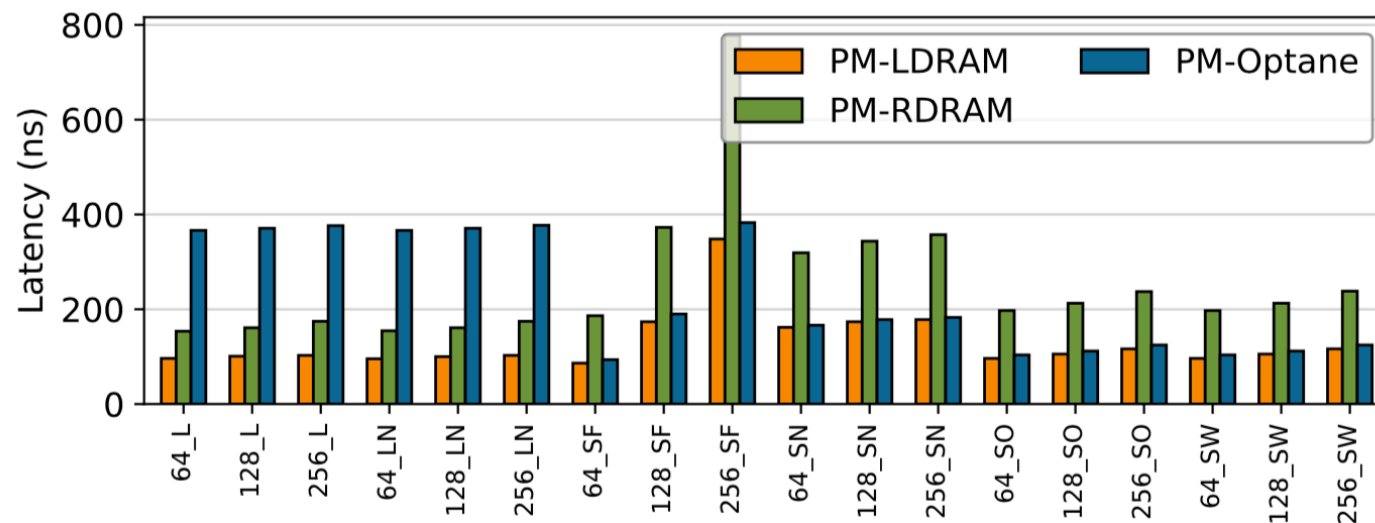
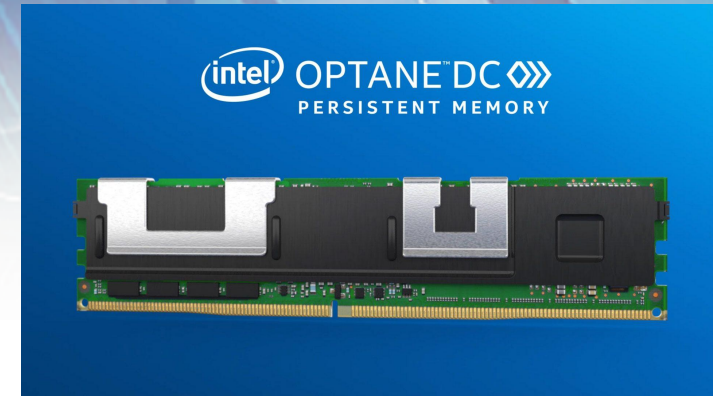


Figure 8: **Memory Instruction Latency** This graph shows the median latency for a variety of ways of accessing persistent memory. Note that for store instructions followed by flushes, there is little performance difference between PM-LDRAM and PM-3DXP, whereas the DRAM outperforms Optane DC memory for load sequences (see data in csvroot/basic/instruction_latency.csv).

Courtesy of NVSL, UCSD arXiv:1903.05714v2

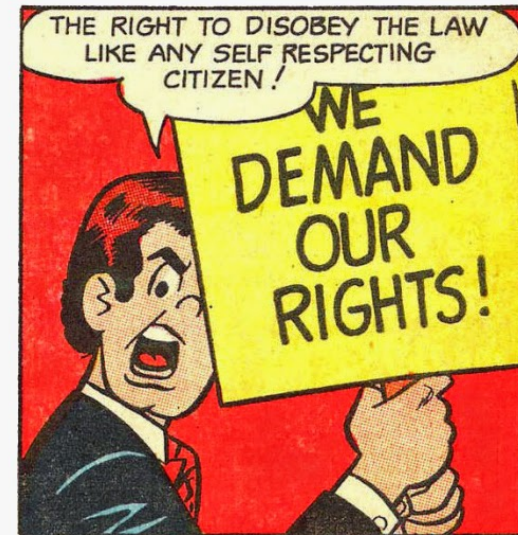
NVRAM



Change the world!



It's Time To Start a Revolution



Peripherals No More!



Change the world!...slowly



It's Time To Start a Revolution



Peripheral



More!



One step at a time...



PAST storage topics of interest?

- **RAID**
 - Increase I/O bandwidth
- **Buffer Caching**
 - Improve latency
- **Swapping**
 - Improve resource sharing



**Revisit &
Rediscover**

Take a fresh look at these old favorites.

PAST storage topics of interest?

- **RAID**
 - Increase I/O bandwidth
- **Buffer Caching**
 - Improve latency
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 - Improve resource sharing



**Revisit &
Rediscover**

Take a fresh look at these old favorites.

NVRAM



SWAN

It's the network, stupid!

NVRAM

Alleviating Garbage Collection Interference through Spatial Separation in All Flash Arrays

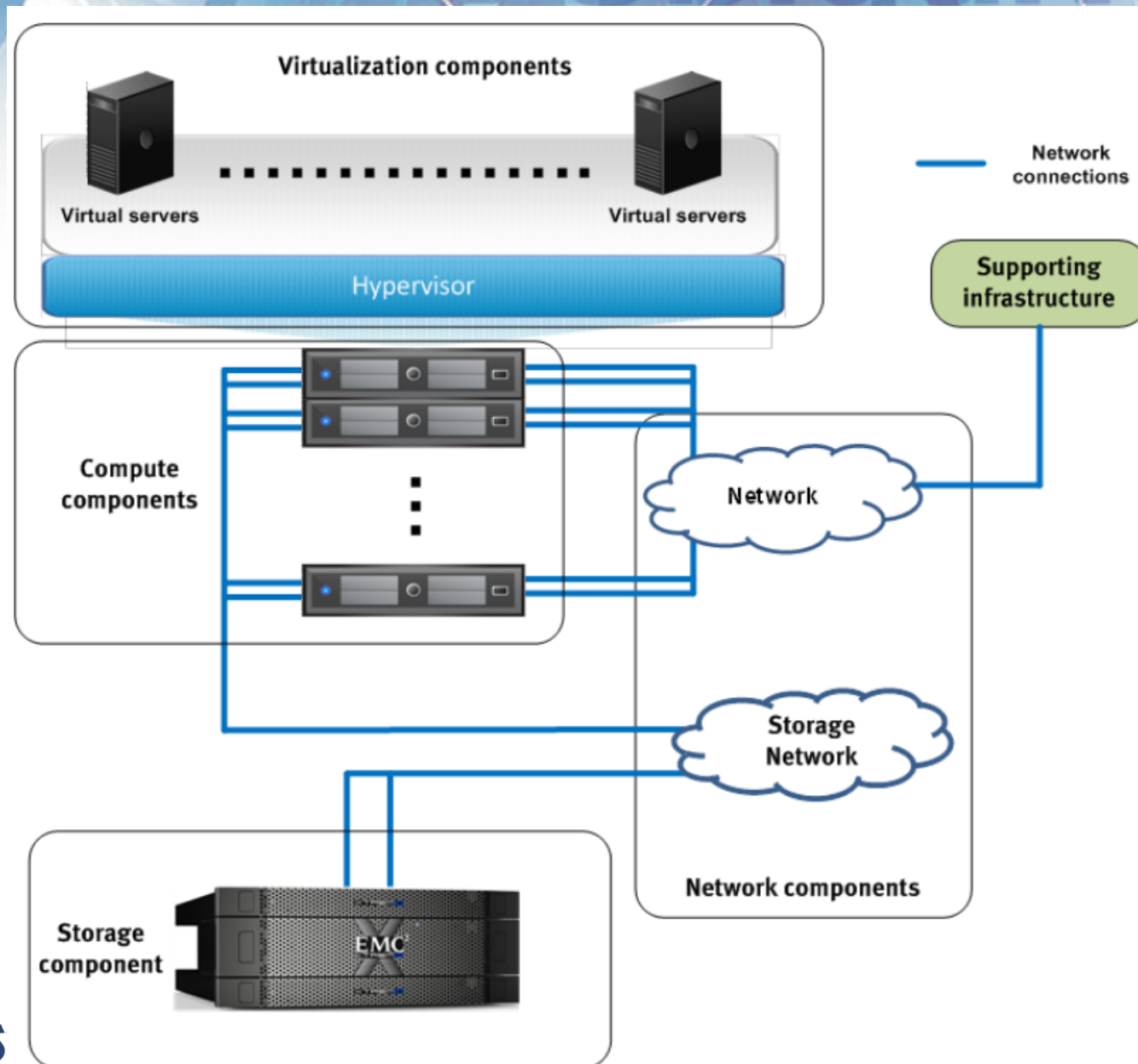
HotStorage '17 & ATC '19

All Flash Array

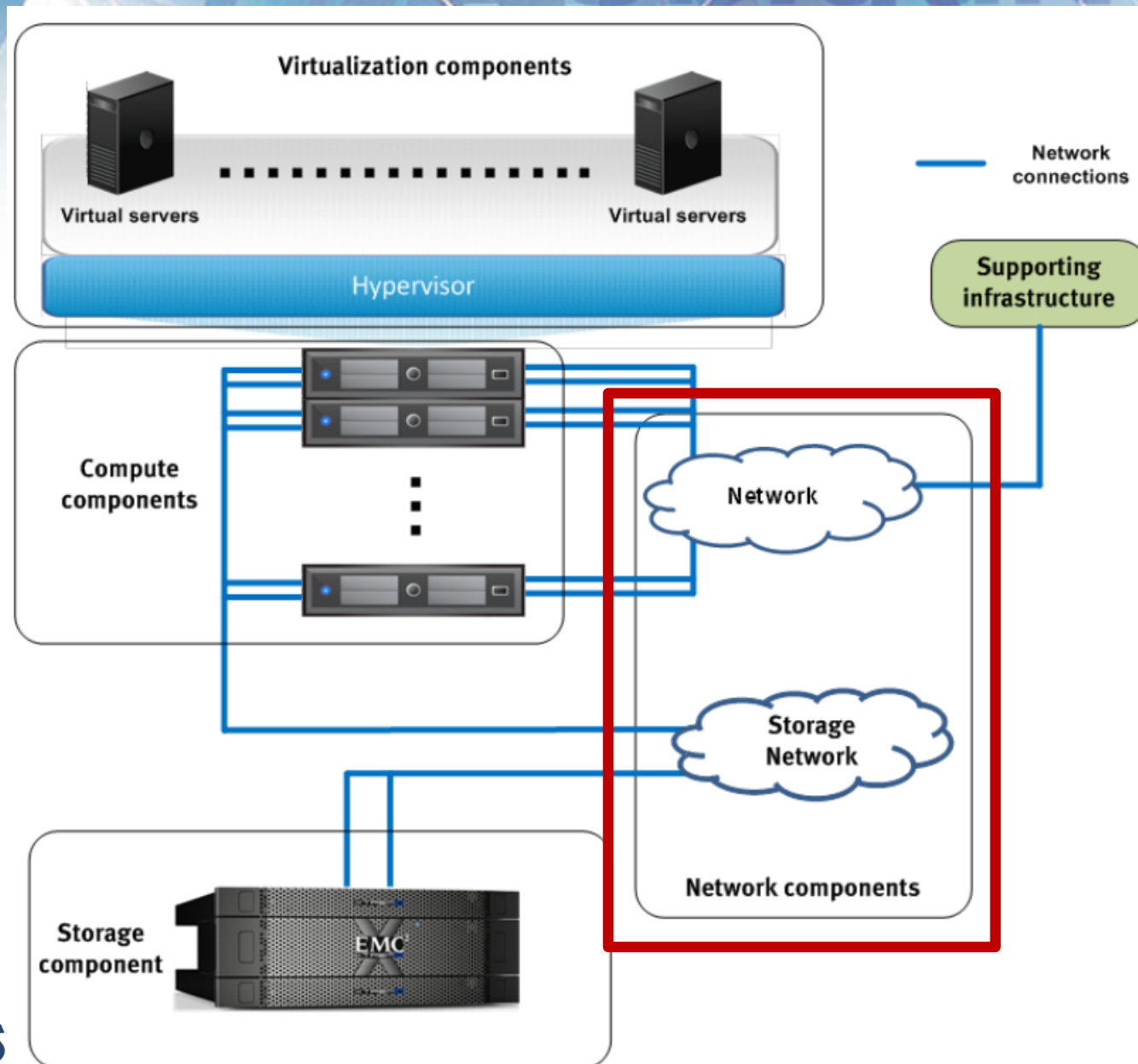
- **All Flash Array (AFA)**
 - Storage infrastructure that contains only flash memory drives
 - Solid-State Array (SSA)



Architecture of All-Flash Array



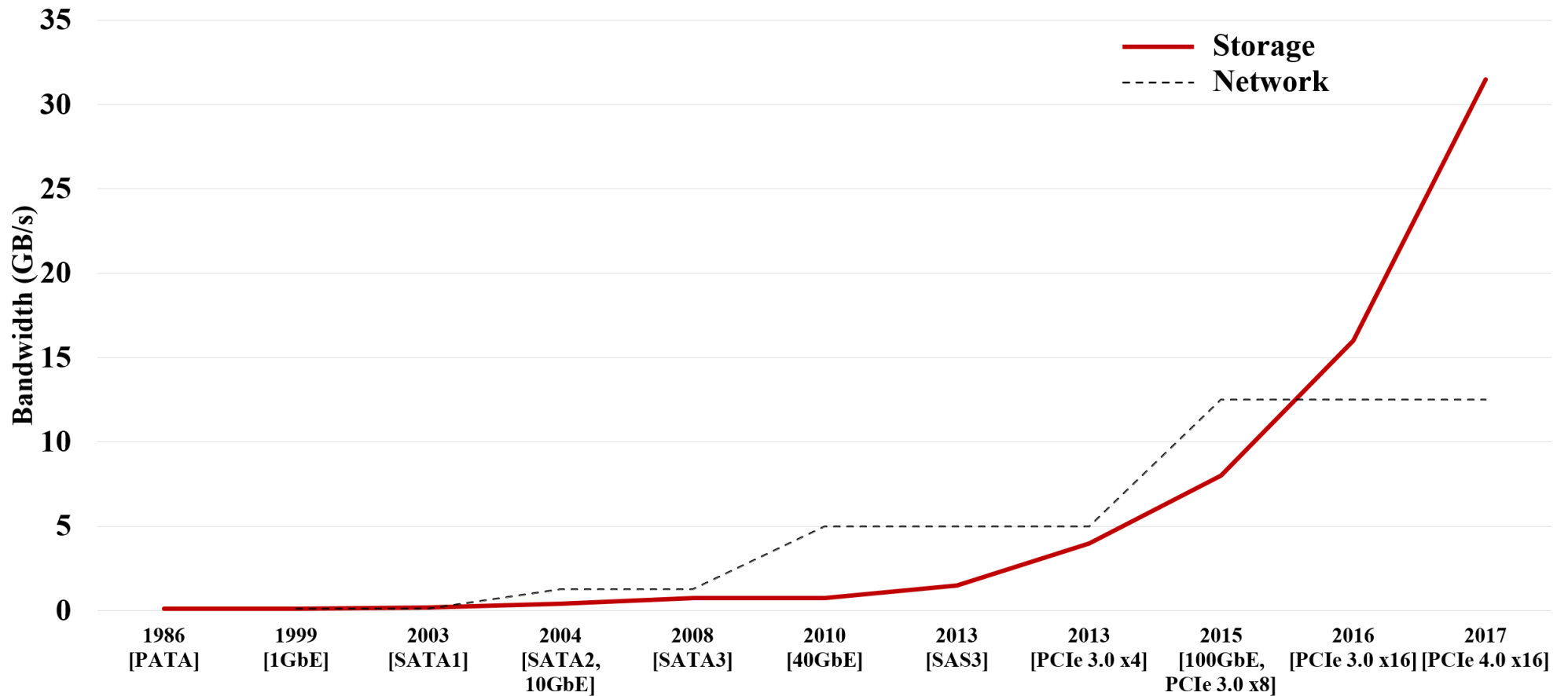
Architecture of All-Flash Array



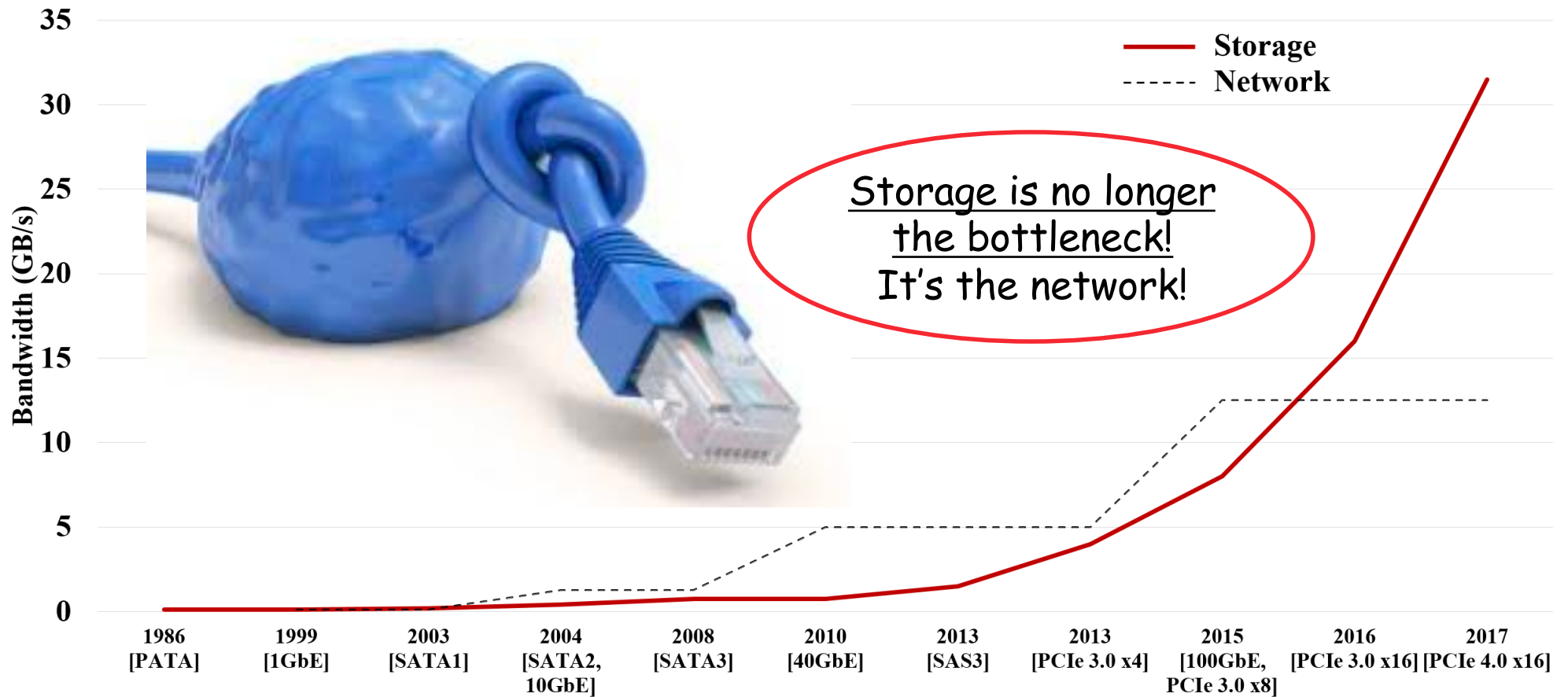
SSD Products for Data Center

Manufacturer	Product Name	Sequential Read/Write (up to GB/s)	Random 4KB Read/Write (up to IOPS)	Interface
Intel	P3700	2.1 / 1	470K / 65K	PCIe 3 * 4
	P3520	1.7 / 1.3	370K / 26K	PCIe 3 * 4
	P3608	5 / 3	850K / 150K	PCIe 3 * 8
	S3710	0.5 / 0.5	85K / 45K	SATA 6Gb/s
Samsung	PM1725a	6.4 / 3	1M / 170K	PCIe 3 * 8
	PM963	2 / 1.2	430K / 40K	PCIe 3 * 4
	PM1633a	1.2 / 0.9	190K / 31K	SAS 3.0
	SM863	0.5 / 0.5	97K / 30K	SATA 6Gb/s

Interface Bandwidth Growth Trend



Interface Bandwidth Growth Trend



Comparison of All-flash Array

NVRAM

	Solid Fire (NetApp)	EMC	Pure Storage	Nimble
Model	SF19210	6X-Brick	M70	AF9000
Capacity	20TB (10 SSDs)	240TB (150 SSDs)	136TB	500TB
Performance (Random I/O)	100K	7GB (900K IOPS * 8KB)	9GB (300K IOPS * 32KB)	350K
Network	20Gb (iSCSI 10Gb * 2port)	240Gb (iSCSI 10Gb * 24port)	40Gb (iSCSI 10Gb * 4port)	40Gb (iSCSI 10Gb * 4port)
Bottleneck	Network	Storage	Network	Network

EMC: <https://www.emc.com/collateral/data-sheet/h12451-xtremio-4-system-specifications-ss.pdf>

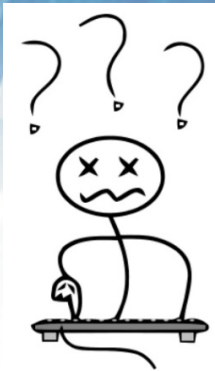
Pure Storage: https://www.purestorage.com/content/dam/purestorage/pdf/datasheets/ps_ds5p_flasharraym_04.pdf

SolidFire: http://info.solidfire.com/rs/solidfire/images/SolidFire_ProductDatasheet.pdf

Nimble storage: <https://www.nimblestorage.com/technology-products/all-flash-array-specifications/>

Comparison of All-flash Array

NVRAM



Do these many SSDs really help?

	Solid Fire (NetApp)	EMC	Pure Storage	Nimble
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EMC: <https://www.emc.com/collateral/data-sheet/h12451-xtremio-4-system-specifications-ss.pdf>

Pure Storage: https://www.purestorage.com/content/dam/purestorage/pdf/datasheets/ps_ds5p_flasharraym_04.pdf

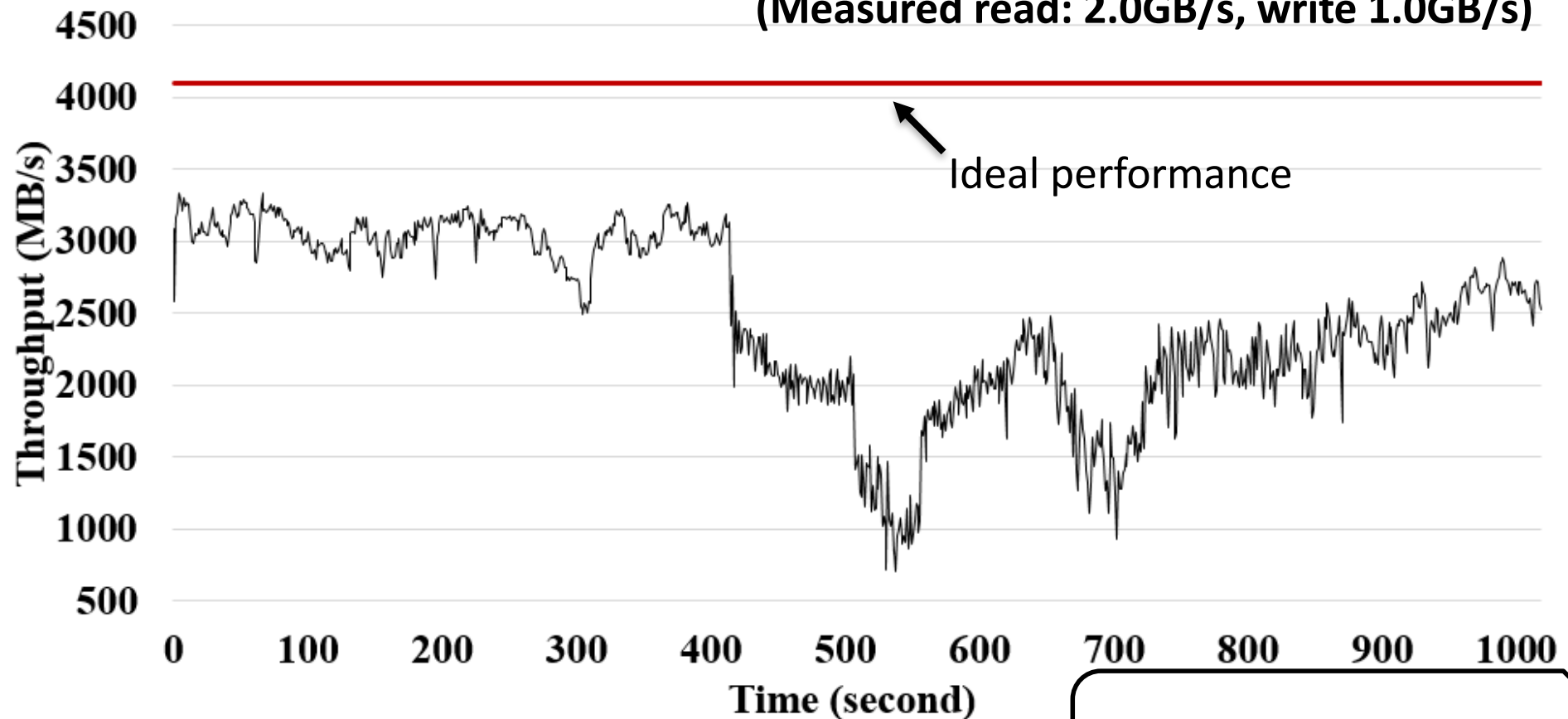
SolidFire: http://info.solidfire.com/rs/solidfire/images/SolidFire_ProductDatasheet.pdf

Nimble storage: <https://www.nimblestorage.com/technology-products/all-flash-array-specifications/>

Experiments with 4 SSD RAID 0

RAID 0 with 4 NVMe SSDs (spec. read: 2.4GB/s, write: 1.2GB/s)

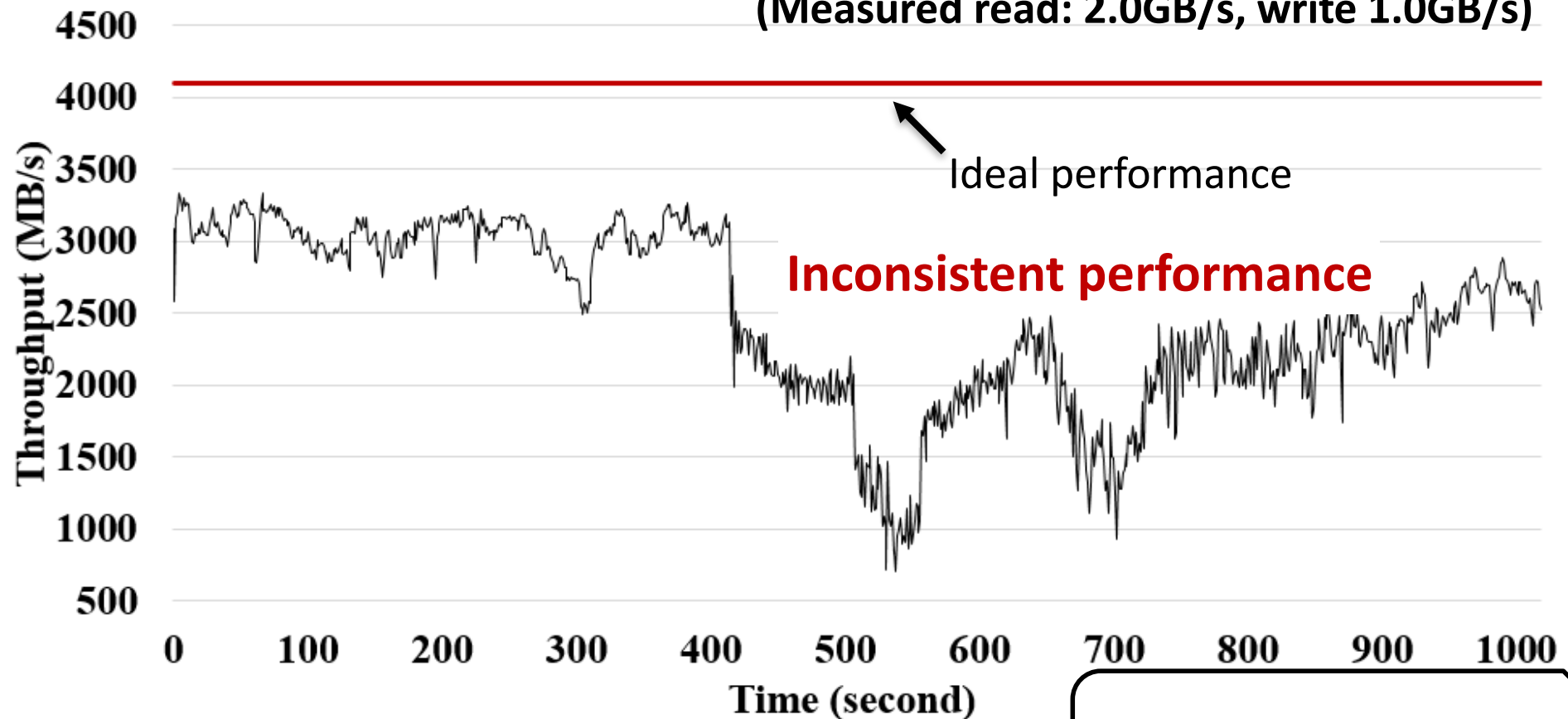
(Measured read: 2.0GB/s, write 1.0GB/s)



Sequential write with
128KB I/O size

Experiments with 4 SSD RAID 0

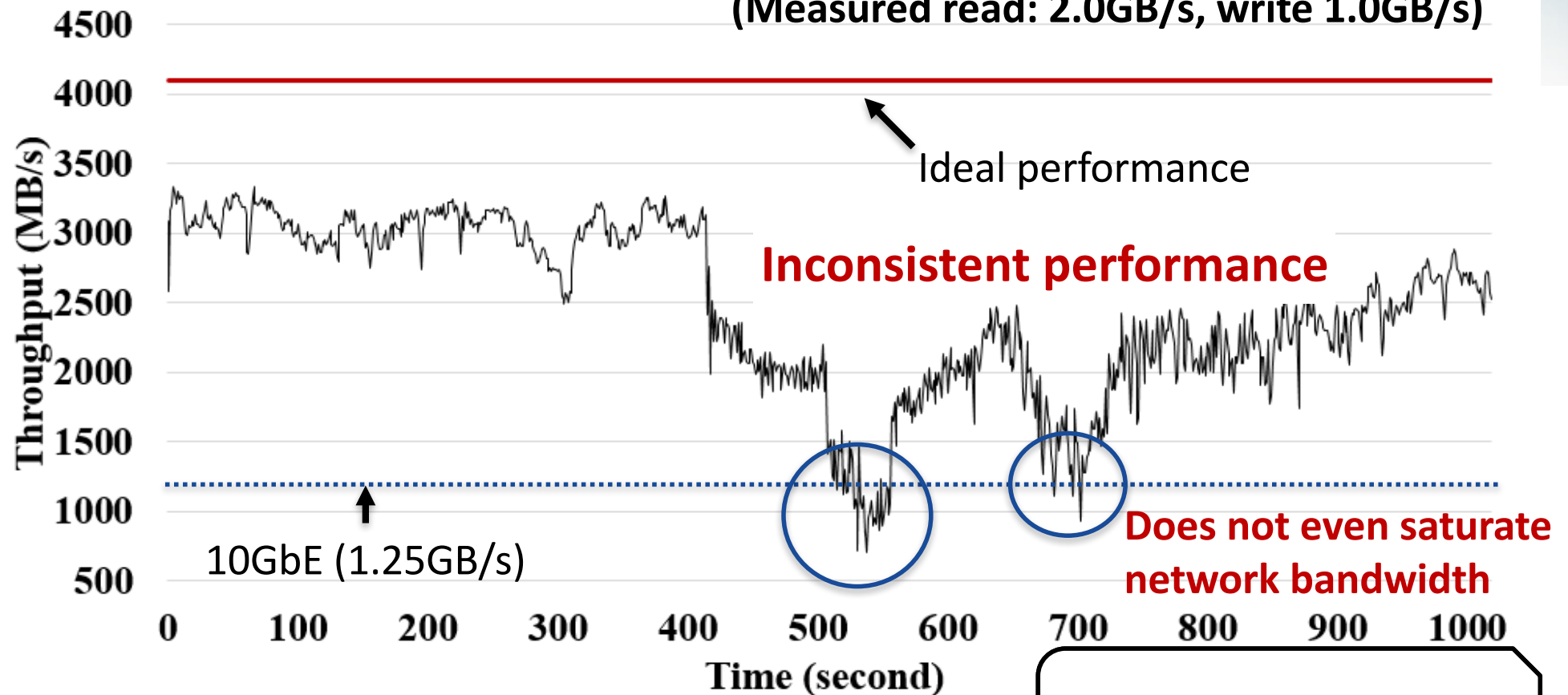
RAID 0 with 4 NVMe SSDs (spec. read: 2.4GB/s, write: 1.2GB/s)
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Experiments with 4 SSD RAID 0

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(Measured read: 2.0GB/s, write 1.0GB/s)



Sequential write with
128KB I/O size

Observations

- Inconsistent performance due to garbage collection
- Performance even limited by network bandwidth

Different approach to arrays of disks

NVRAM

- Inconsistent performance due to garbage collection

Get rid of garbage collection!

- Performance even limited by network bandwidth

Provide full network performance!





Our goal

Sustained, consistent
full network bandwidth performance!

Design of SWAN

- **Our system**
 - SWAN (**S**patial separation **W**ithin an **A**rray of **SSD**s on a **N**etwork)
- **Goals**
 - Provide sustainable high performance for AFA
 - Alleviating GC interference at both SSD-level and AFA-level
- **Approach**
 - Spatial separation of application I/O and AFA I/O
 - Minimize GC interference by organizing SSDs into two-dimensional array

Comparison of RAID schemes

 write req.
  read req.

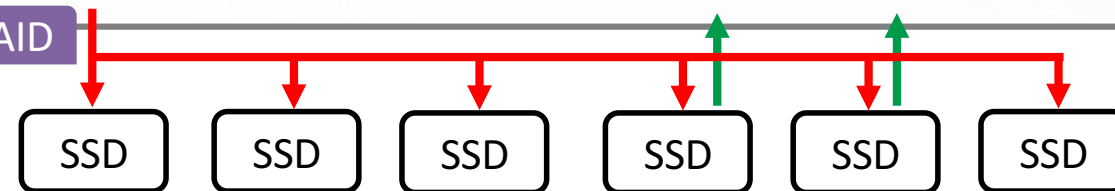
Traditional RAID

RAID



Log-structured writing on RAID

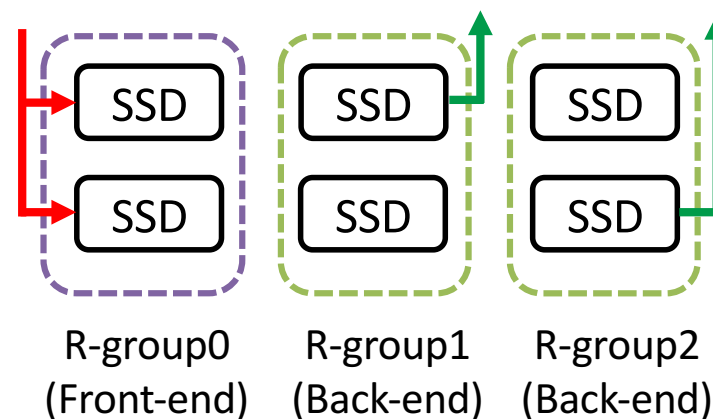
Log-RAID



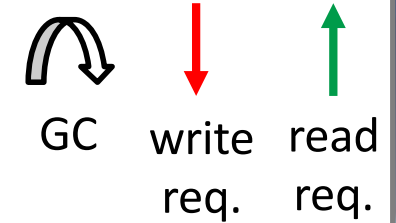
SWAN

- Two dimensional array
- Log-structured writing per R-group
- Front-end servers write requests
- Back-end is used for AFA-level GC

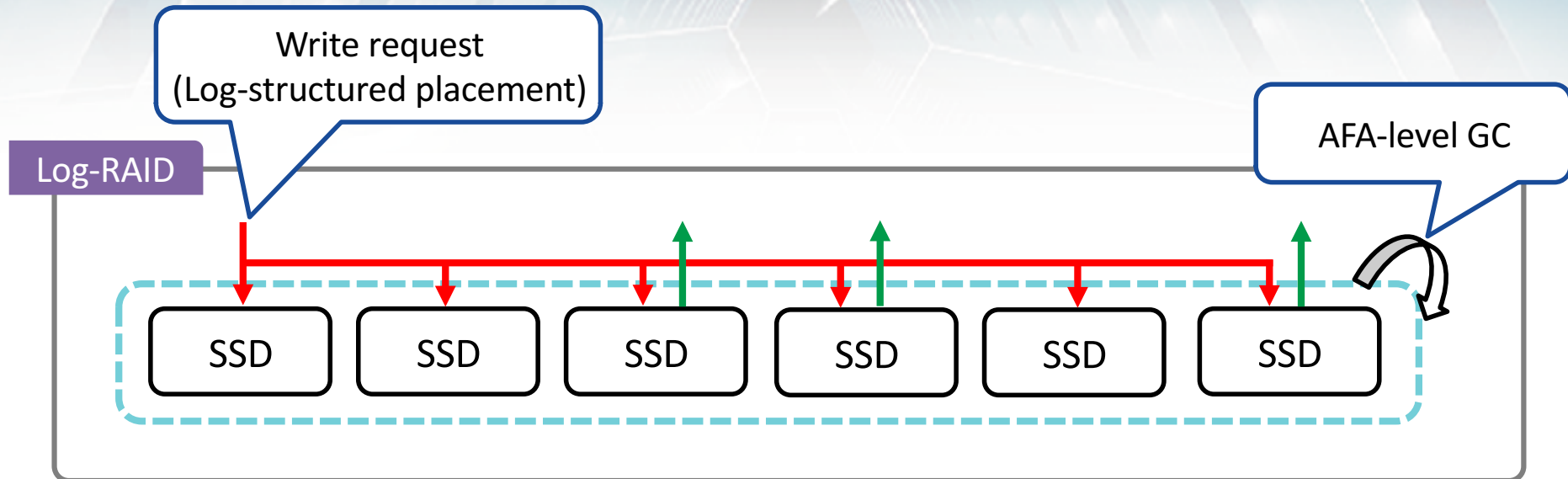
SWAN



How Log-RAID Works



Key operations of Log-RAID



[9] CHIUUEH, T.-C., TSAO, W., SUN, H.-C., CHIEN, T.-F., CHANG, A.-N., AND CHEN, C.-D. Software orchestrated flash array. In *Proceedings of International Conference on Systems and Storage (SYSTOR)* (2014), pp. 14:1–14:11.

[21] IOANNOU, N., KOURTIS, K., AND KOLTSIDAS, I. Elevating commodity storage with the SALSA host translation layer. In *Proceedings of the 26th IEEE International Symposium on Modeling, Analysis, and Simulation of Computer and Telecommunication Systems (MAS-COTS)* (2018), pp. 277–292.

[10] COLGROVE, J., DAVIS, J. D., HAYES, J., MILLER, E. L., SANDVIG, C., SEARS, R., TAMCHES, A., VACHHARAJANI, N., AND WANG, F. Purity: Building Fast, Highly-Available Enterprise Flash Storage from Commodity Components. In *Proceedings of the ACM SIGMOD International Conference on Management of Data* (2015), pp. 1683–1694.

How SWAN Works



GC



TRIM

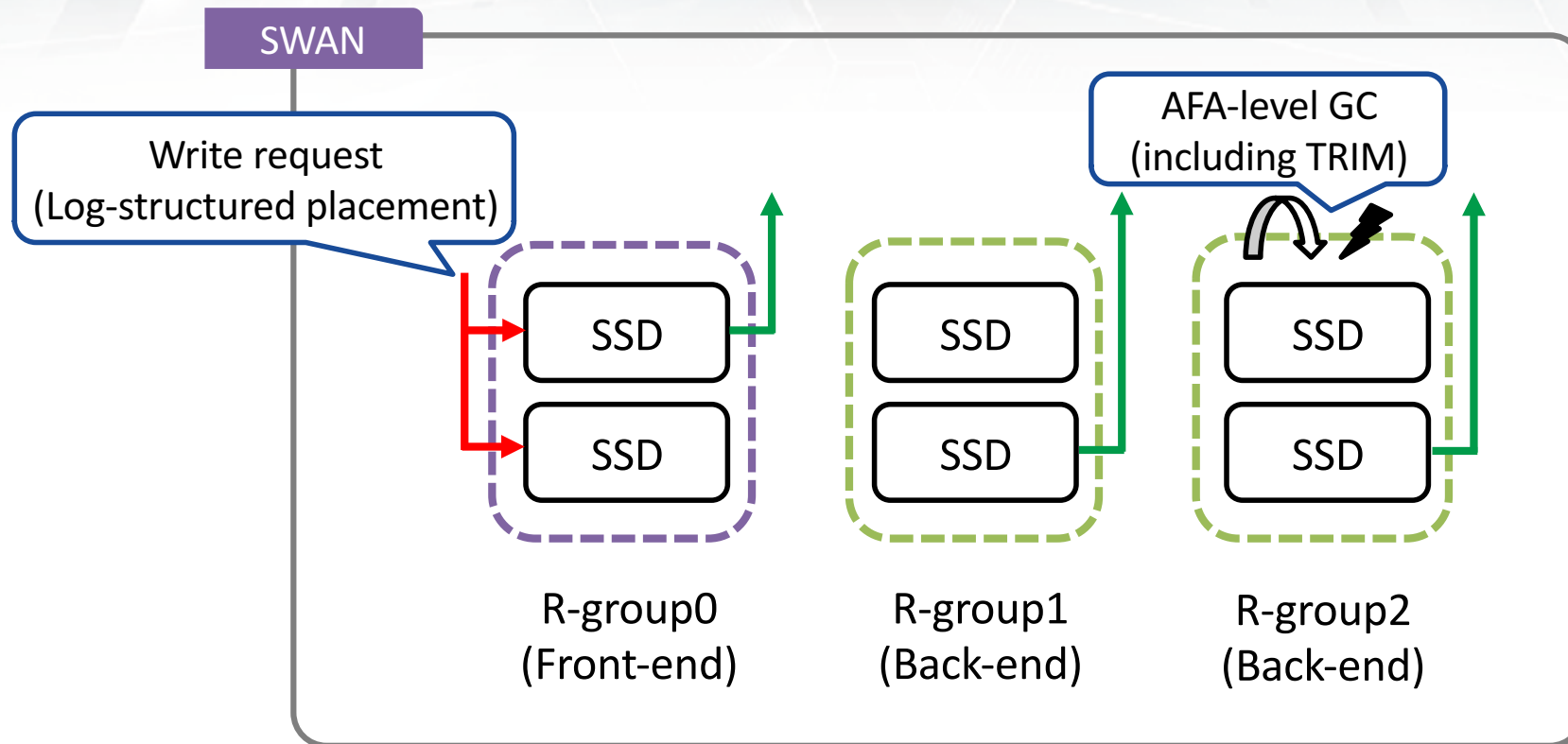


write
req.

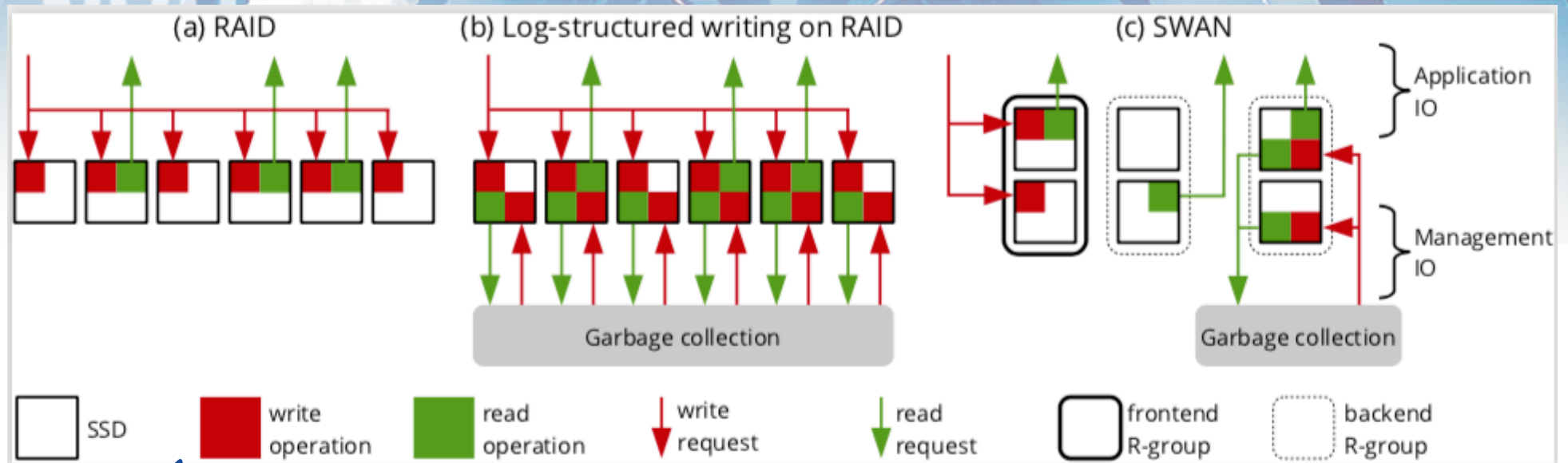


read
req.

Key operations of SWAN



I/O operation in All Flash Array



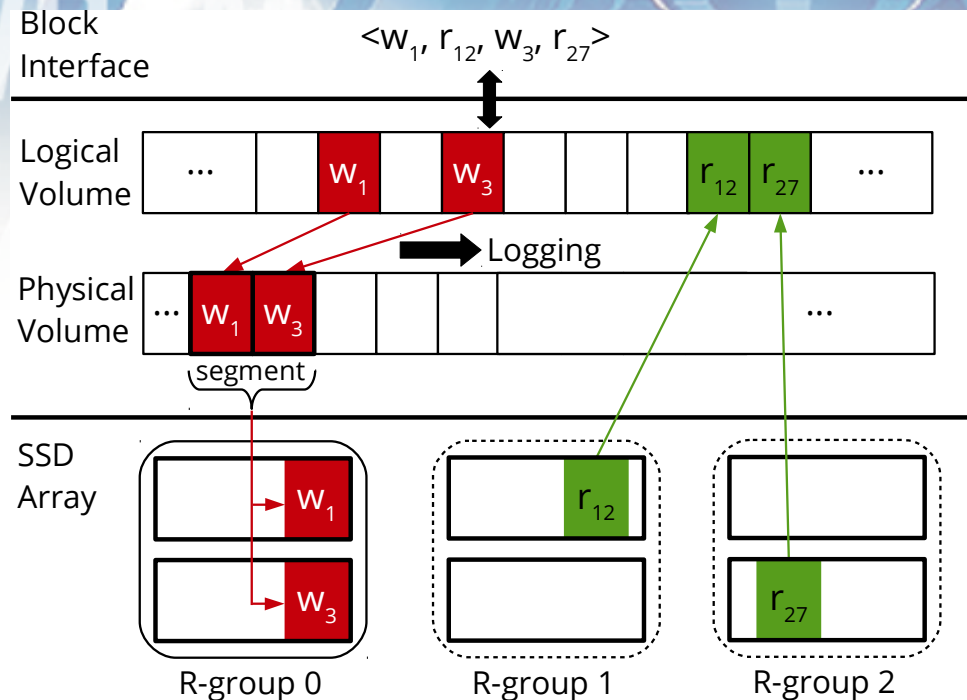
Susceptible to performance degradation due to high GC overhead inside SSD (Due to random writes)

AFA-level GC I/O may significantly interfere with application I/O

Spatial separation of application I/O and AFA-level GC I/O to minimize I/O interference

Handling Read/Write Req. in SWAN

w_n : write req. for block n
 r_n : read req. for block n



Conf.

- R-group 0: Front-end
- R-group 1,2: Back-end
- Read/write req. arrives via block interface

■ Operations

- SWAN appends write req. to the log and issues write req. to the front-end
- Read req. will be served by any R-group holding the requested blocks

Evaluation

■ Environment

- Dell R730 server equipped with 2 Xeon CPUs and 64GB DRAM
- Samsung 850 PRO 128GB * 9

■ Target config.

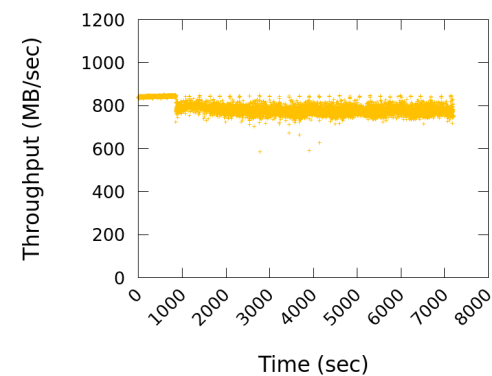
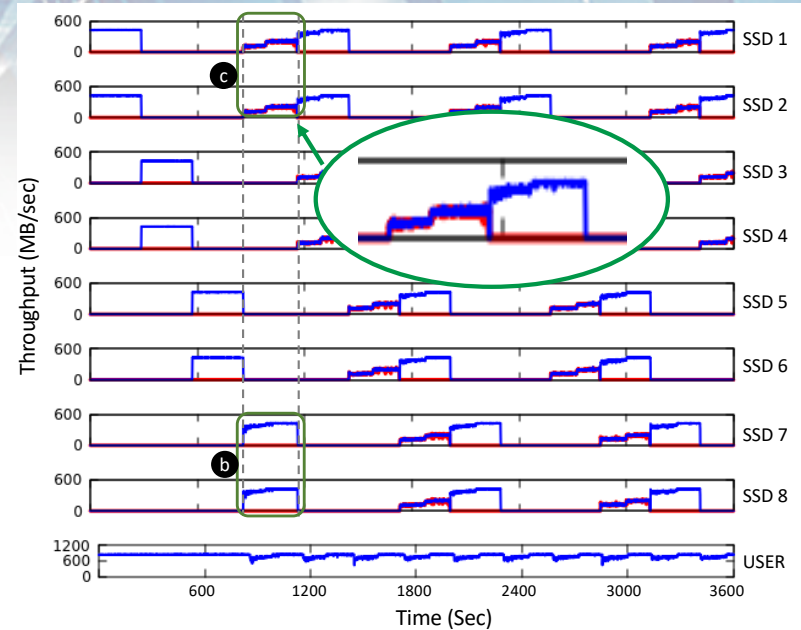
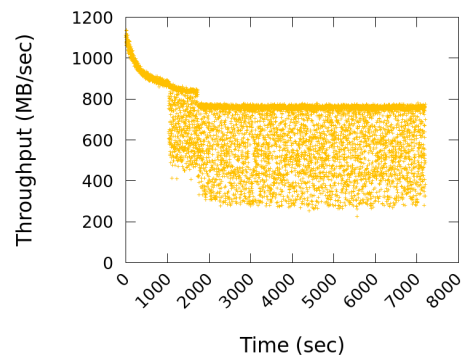
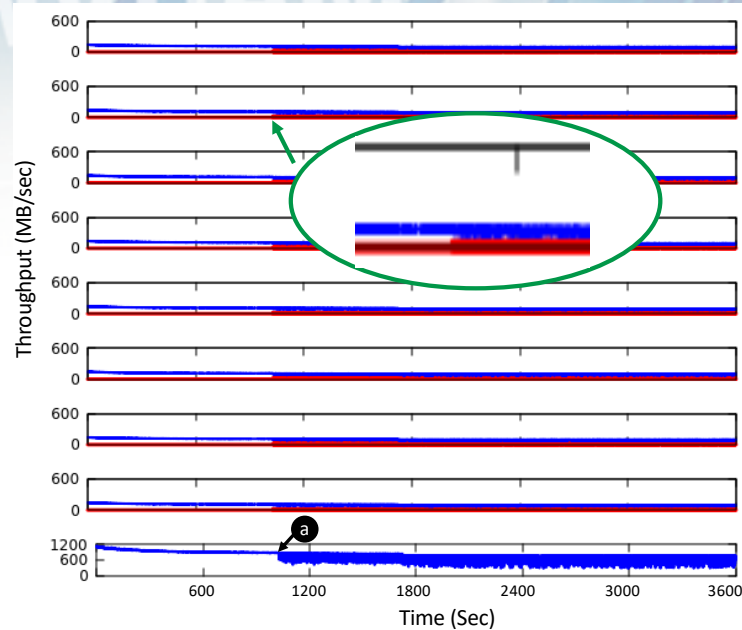
- RAID-0/4/5
- Log-RAID-0/4
- SWAN-0/4

■ Workloads

- Microbenchmark
- YCSB-A, B, C, and D

Analysis of GC Behavior

Random write workload



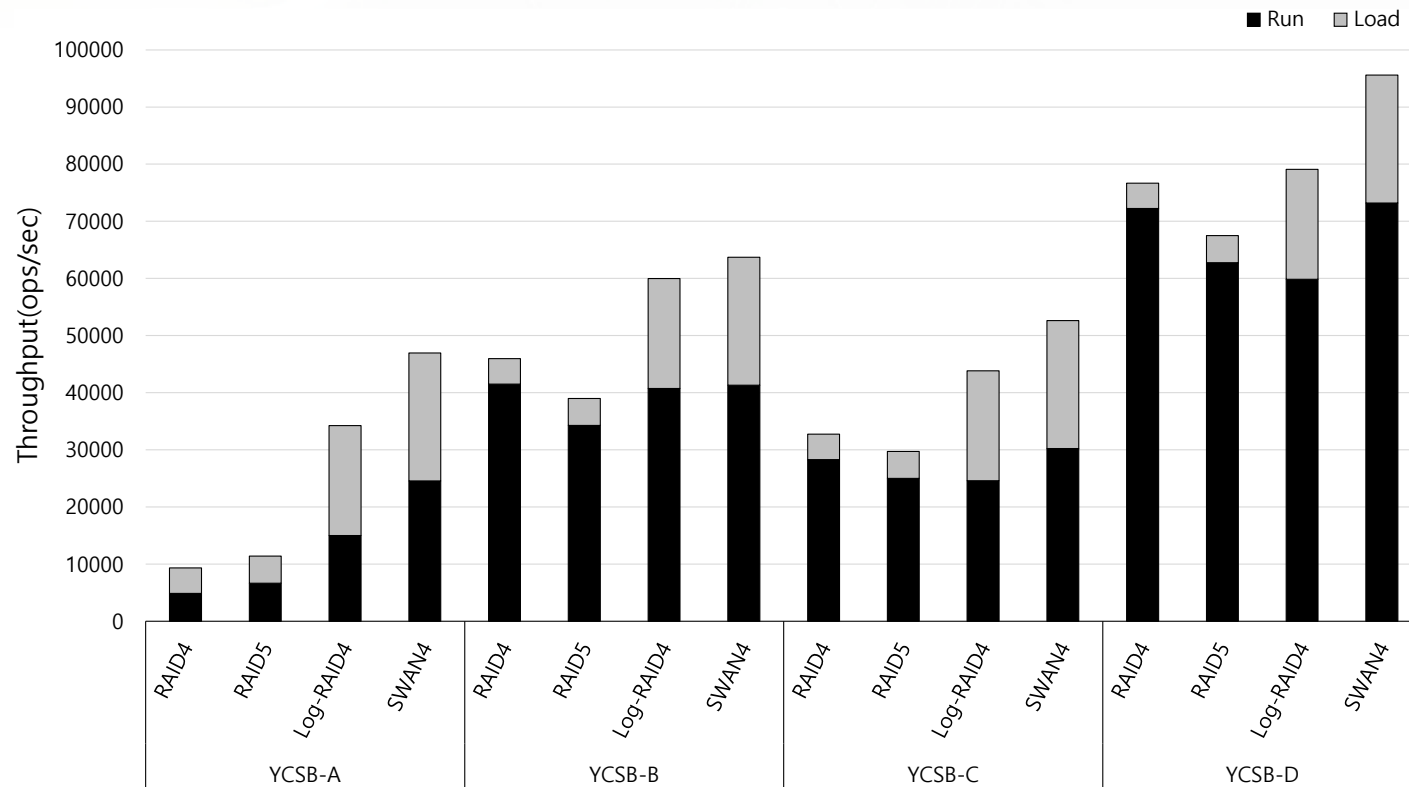
Log-RAID (8 SSDs)

SWAN (4 R-groups / 2 SSD per R-group)

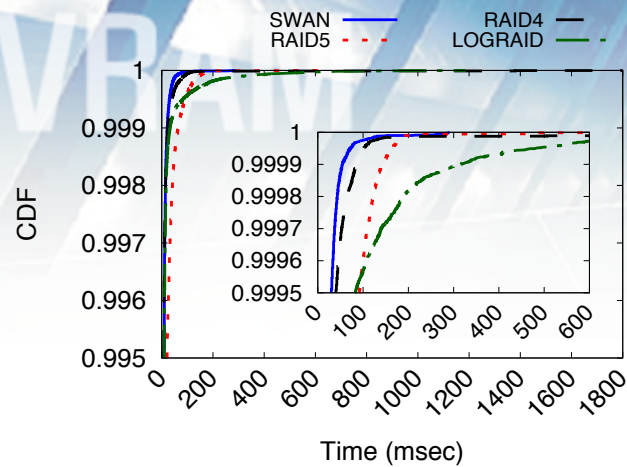
Throughput Results

■ Configuration

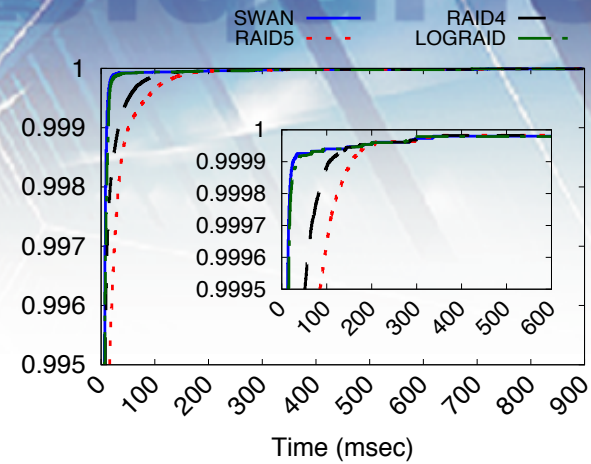
- RAID4/5: 8 data SSDs + 1 parity SSD
- Log-RAID: 8 data SSDs + 1 parity SSD
- SWAN4: 3 R-group with 2 data SSDs and 1 parity SSD per R-group



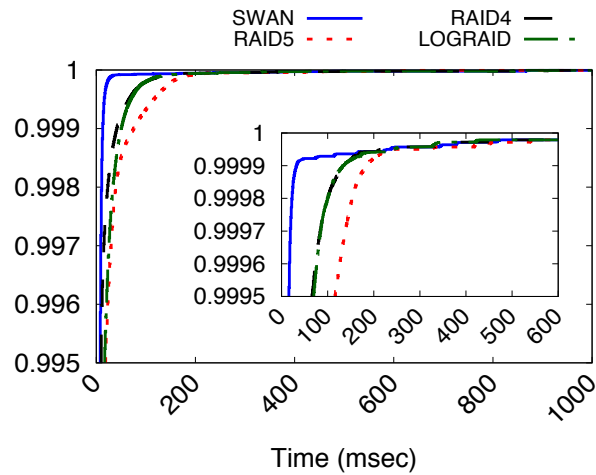
Read Latency Results (CDF)



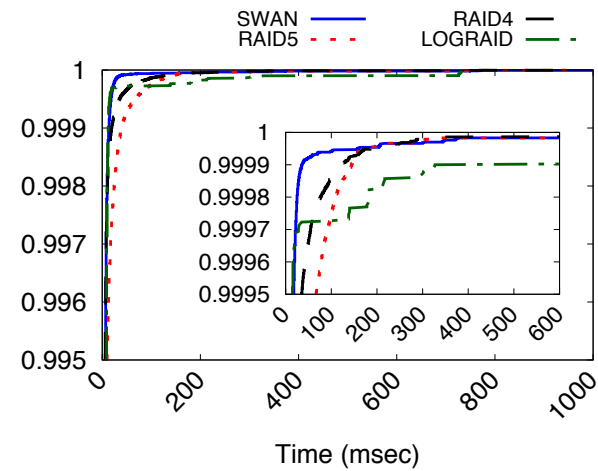
YCSB-A



YCSB-B



YCSB-C



YCSB-D

Summary

- **Proposed SWAN**
 - New management policy for All Flash Array
- **Key idea of SWAN**
 - Decouple GC I/Os from normal ones by partitioning the SSD array into 2 groups
- **Benefits of Swan**
 - SSD can be simpler

It's the network, stupid!

PAST storage topics of interest?

- **RAID**
 - Increase I/O bandwidth
- **Buffer Caching**
 - Improve latency
- **Swapping**
 - Improve resource sharing



Revisit & Rediscover

Take a fresh look at these old favorites.

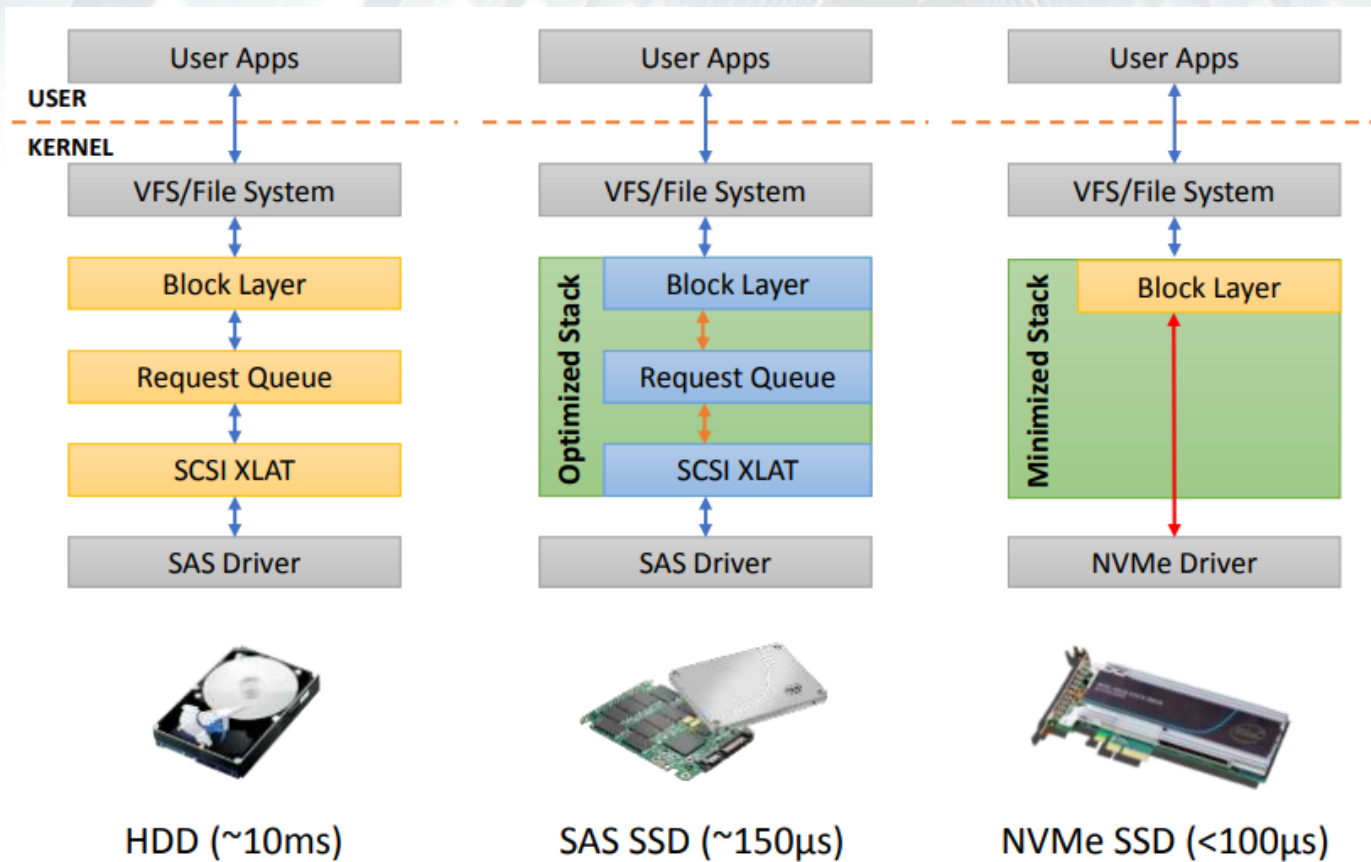
NVRAM



First Responder

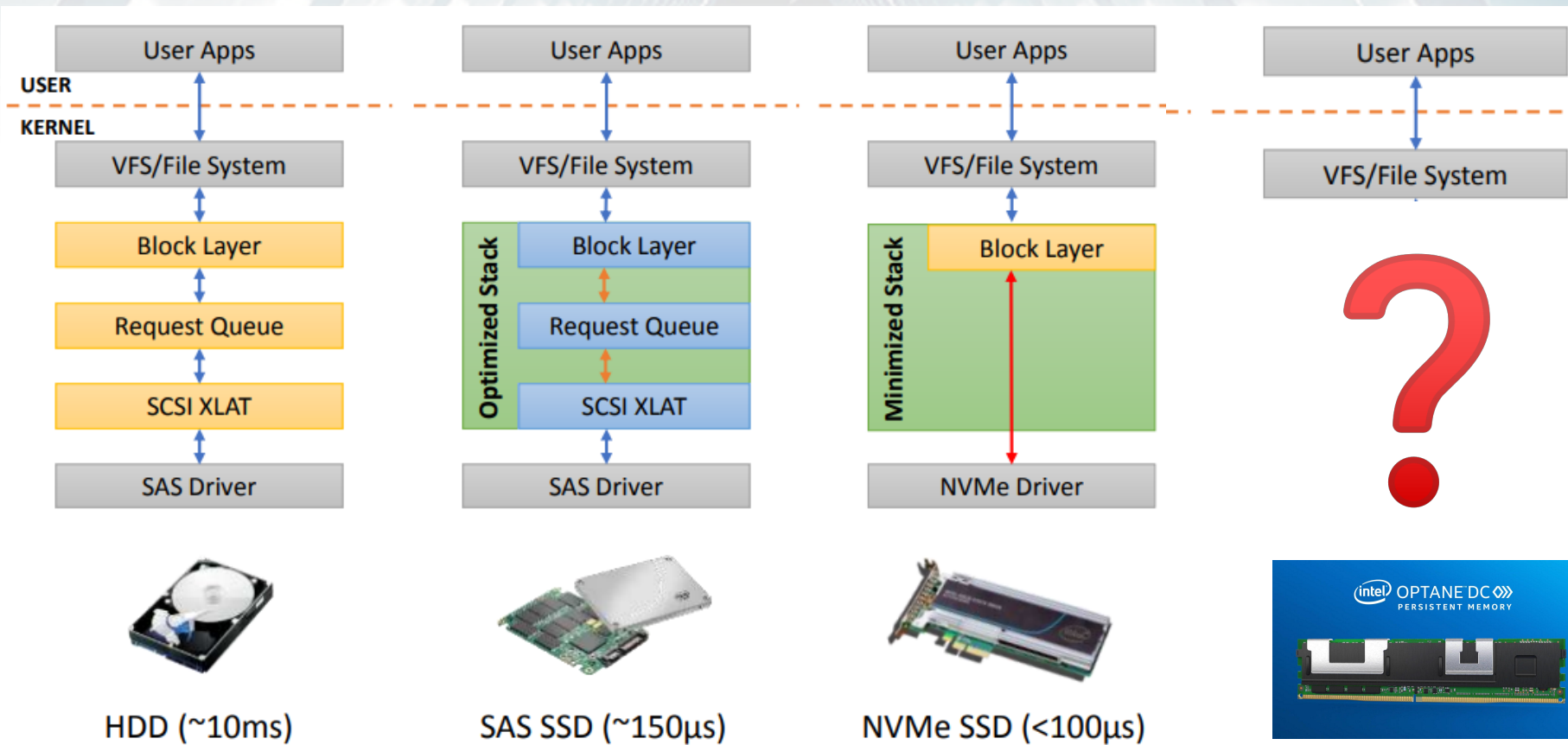
It's the storage stack, stupid!

Evolution of storage stack



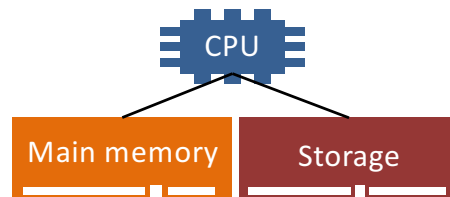
Evolution of storage stack

NV RAM

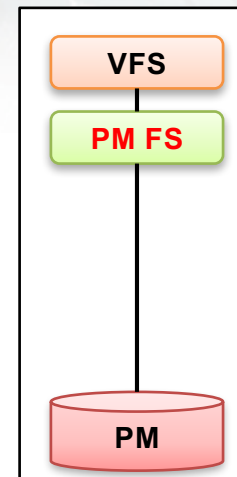


PM Targeted File Systems

- Designed to reap PM performance



PM as Storage



PM-aware File System

SOSP 2009

SC 2011

EuroSys 2014

EuroSys 2014

EuroSys 2016

FAST '16, SOSP '17

SOSP 2017

“BPFS (Better I/O Through Byte-Addressable, Persistent Memory)”

“SCMFS (SCMFS: A File System for Storage Class Memory)”

“PMFS (System Software for Persistent Memory)”

“Aerie (Aerie: Flexible File-System Interfaces to Storage-Class Memory)”

“HiNFS (A High Performance File System for Non-Volatile Main Memory)”

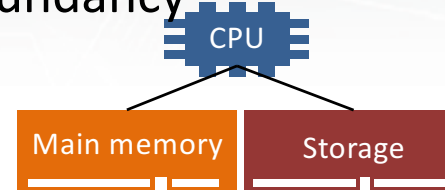
“NOVA (NOVA-Fortis: A Fault-Tolerant Non-Volatile Main Memory File System)”

“Strata (Strata: A Cross Media File System)”

BUT...

■ DAX approach

- Weak reliability, data integrity, redundancy
- PM as end destination media



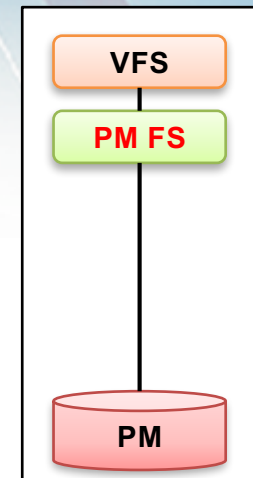
PM as Storage

■ PM only

- Replace traditional storage?
- Exception: Strata and Ziggurat

■ Lengthy process to maturity

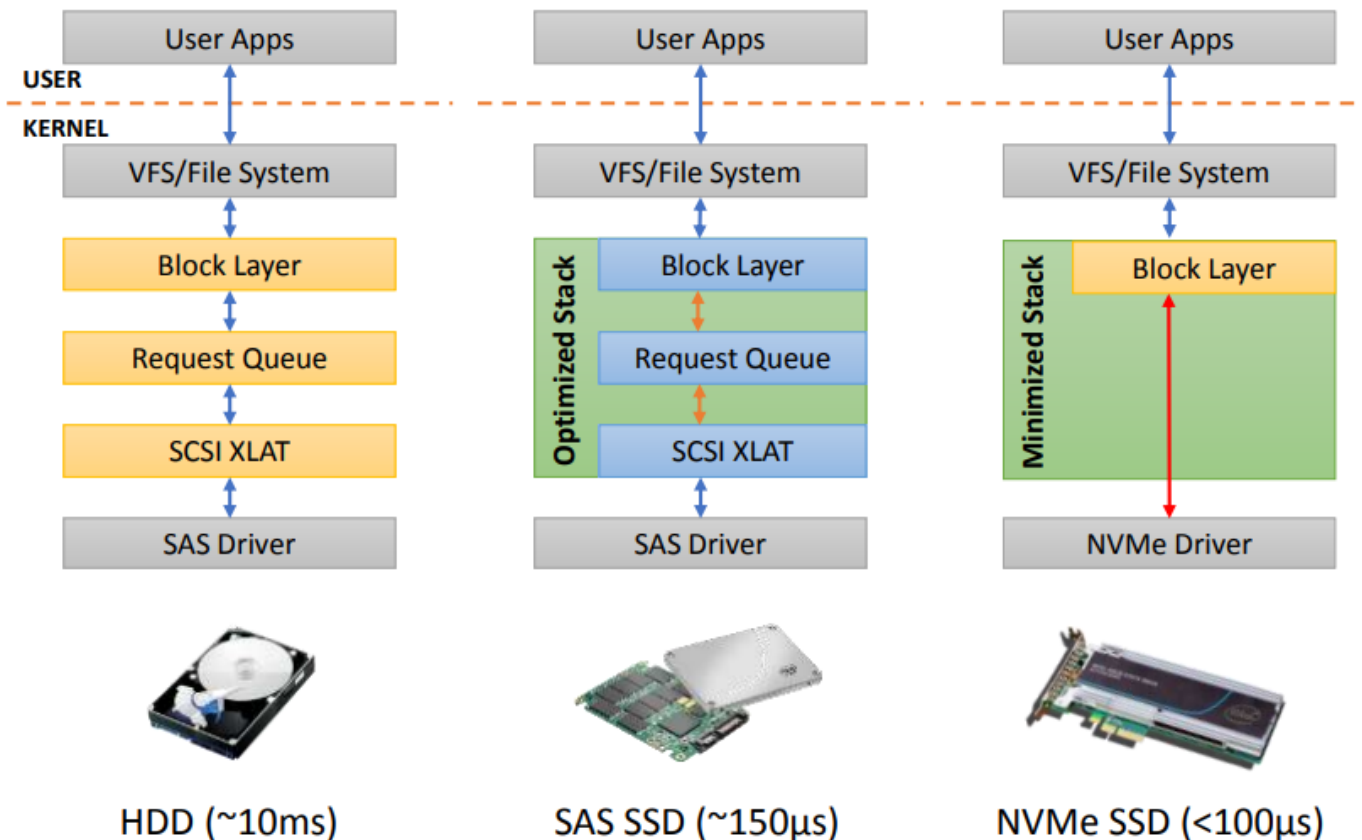
- Ext4...still in progress
- Wisdom with age



PM-aware File System

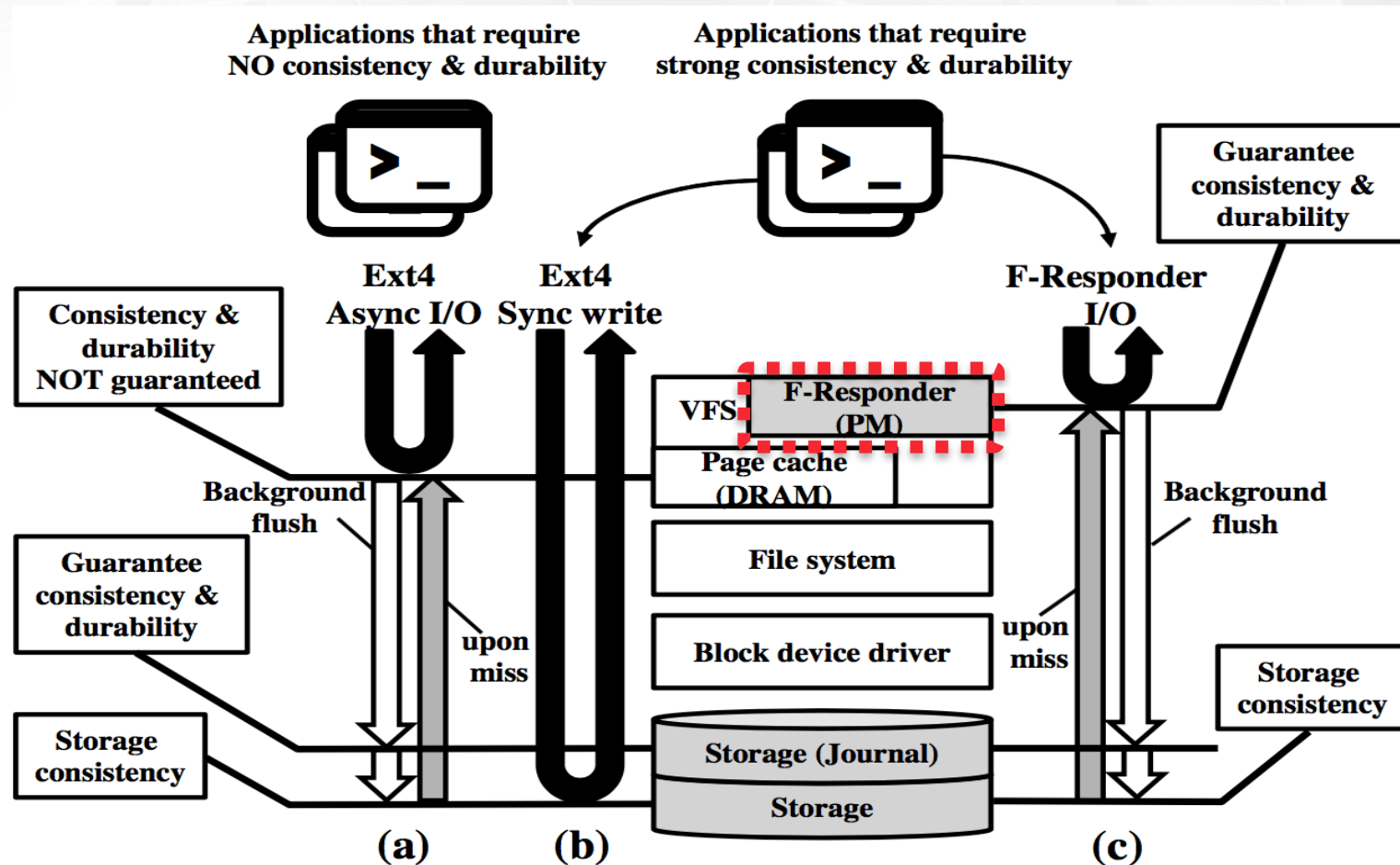
Our Goal

- Keep legacy file system and storage media "as-is"
- Integrate PM for performance and durability/consistency



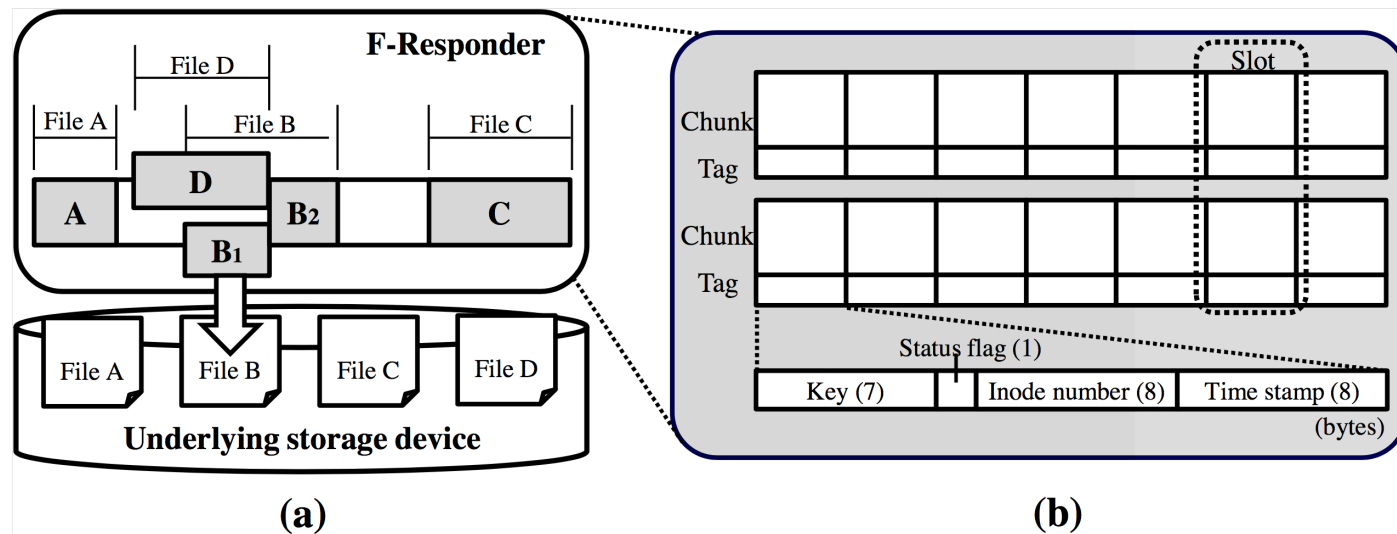
First Responder

- Overall architecture



Design

- Static placement in “buffer cache”
- Sufficient large “cache”
 - Replacement policy (almost) agnostic
- Background flush to underlying storage device
 - Hide storage stack overhead



Performance evaluation

■ System configuration and benchmarks

Table 1. System configuration

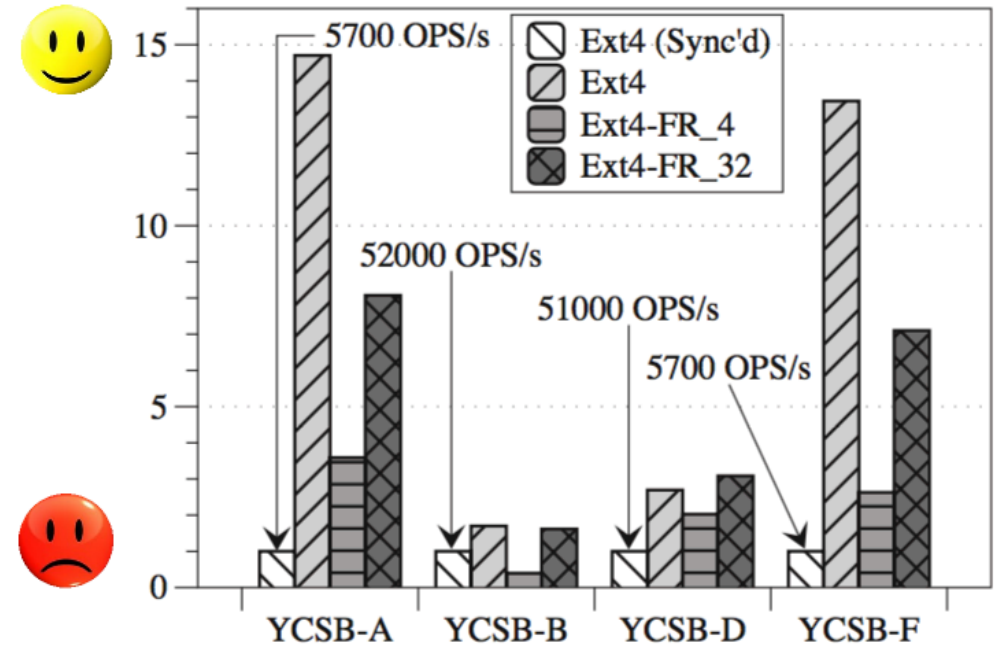
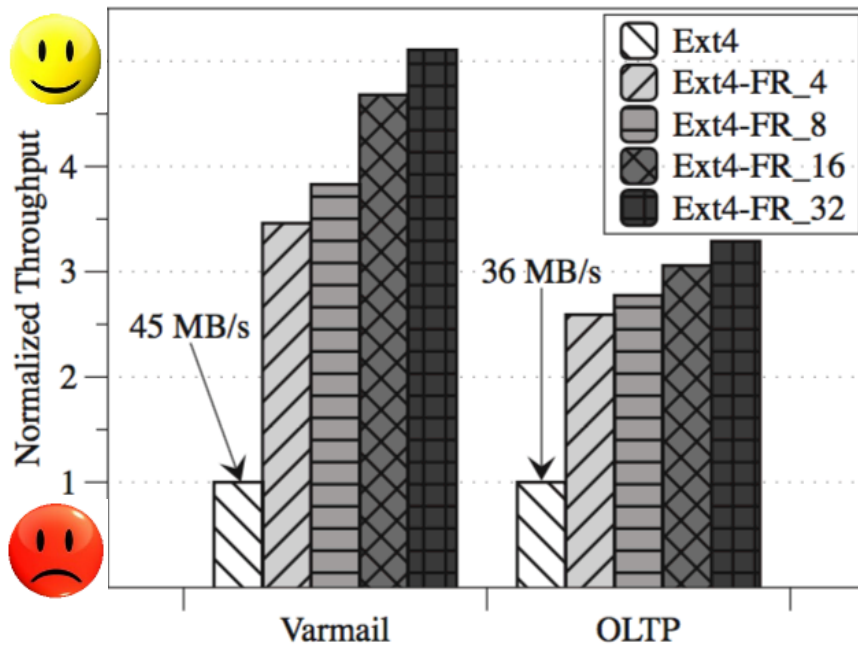
	Description
CPU	Intel Xeon E5-2620V3 (6 cores / 12 threads) × 2
Memory	Samsung DDR4 16GB PC4-17000 × 16 (256GB)
Storage	Samsung V-NAND SSD 850 PRO 256GB
OS	Linux Ubuntu 16.04 LTS (64bit) kernel v4.18

Table 2. Characteristics of benchmarks

Filebench	R:W	Mean file size	# of files	# of threads
Varmail	1:1	32KB	800K	50
OLTP	1:1	1.5GB	20	W:10 R:200
Key-value store	R:W	Record selection	Dataset size	# of threads
YCSB-A	1:1	Zipfian	12GB	20
YCSB-B	19:1	Zipfian	12GB	20
YCSB-D	19:1	Latest	12GB	20
YCSB-F	1:1	Zipfian	12GB	20

Performance evaluation

Overall performance

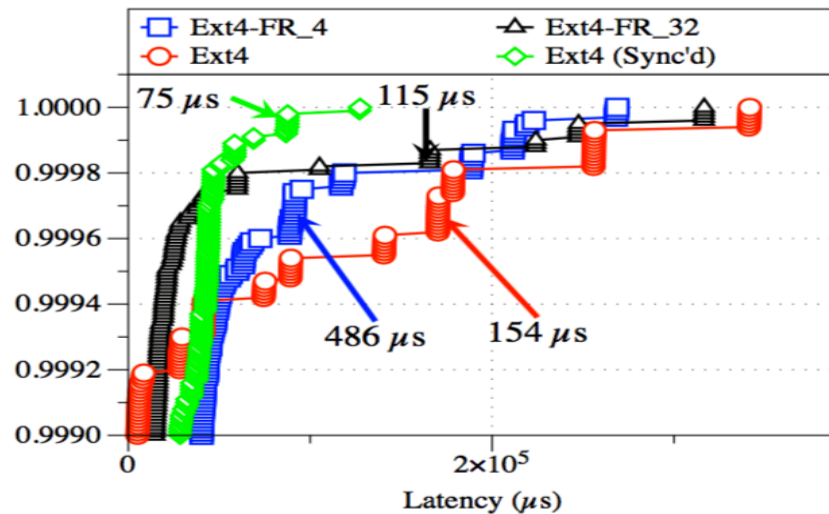


(a) Varmail and OLTP performance relative to Ext4 (async)

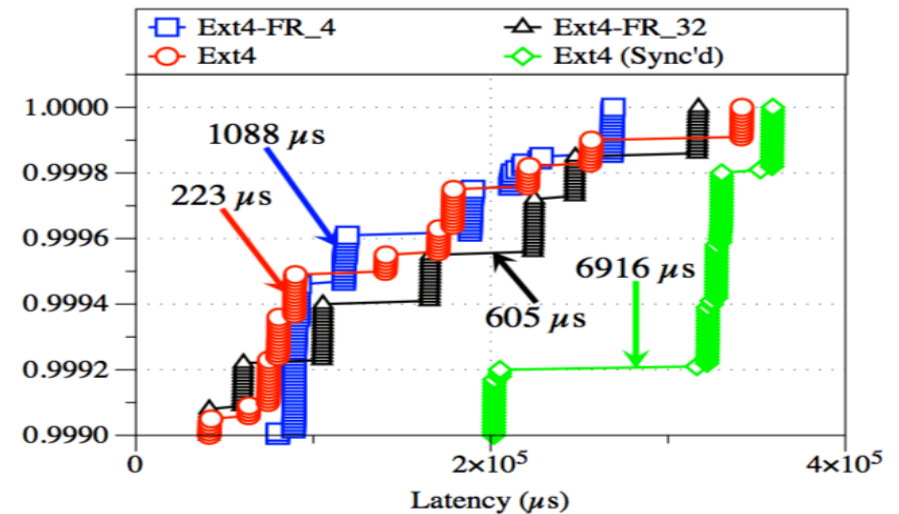
(b) YCSB (with sync mode RocksDB) performance relative to Ext4

Performance evaluation

YCSB-A latency results



(a) Read

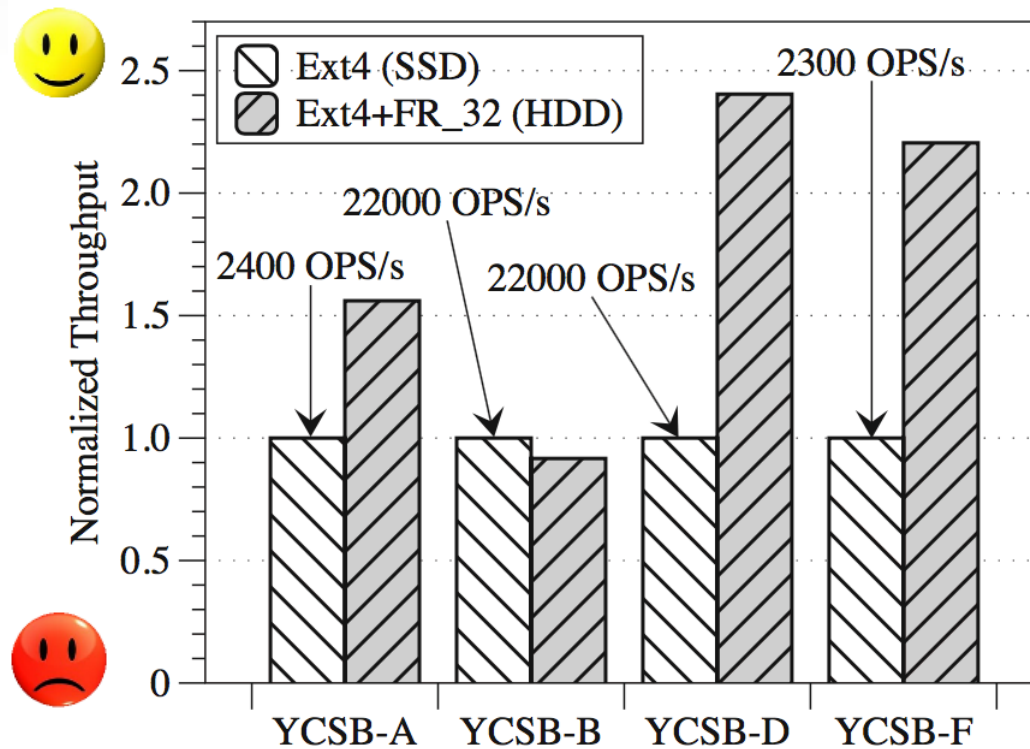


(b) Update

- ✦ In F-Responder, consistency and durability can be guaranteed without much loss in performance
- ✦ Sync mode reads the average is smallest and the tail is very short
- ✦ F-Responder-32GB does better than sync mode on Ext4, but worse than async mode on Ext4

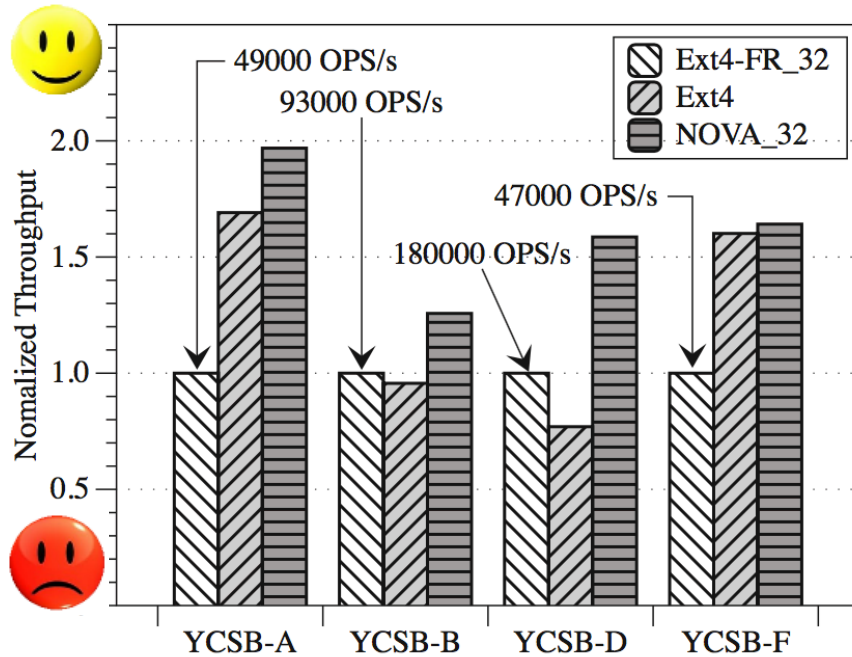
Performance evaluation

- F-Responder with HDD

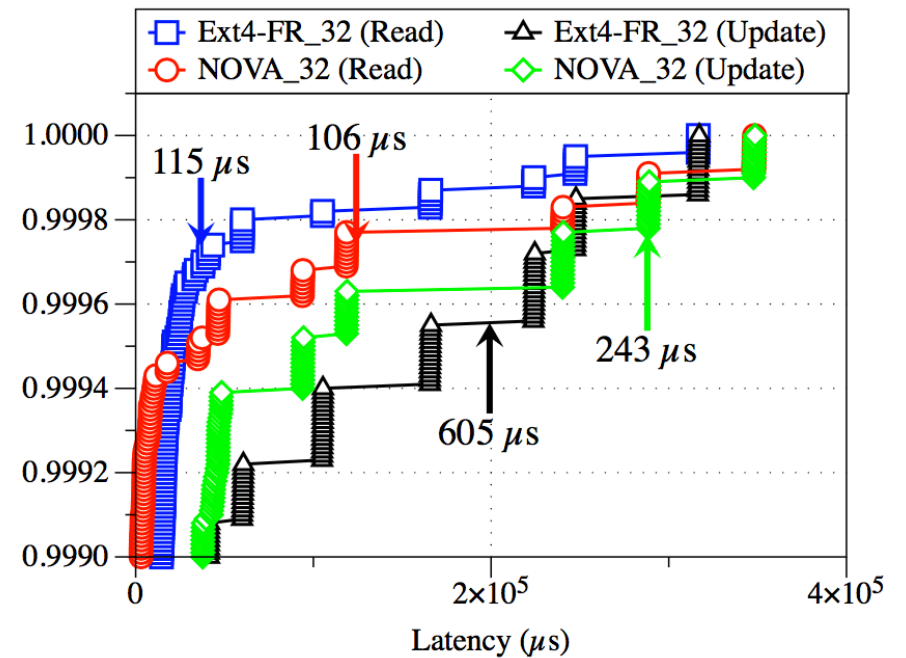


Performance evaluation

Comparison to NOVA-Fortis



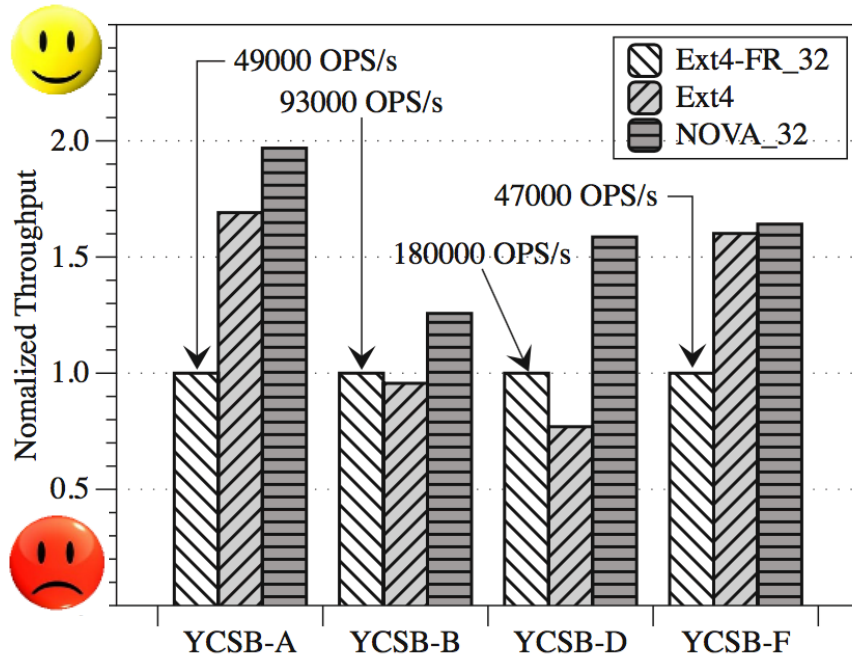
(a) YCSB Workloads



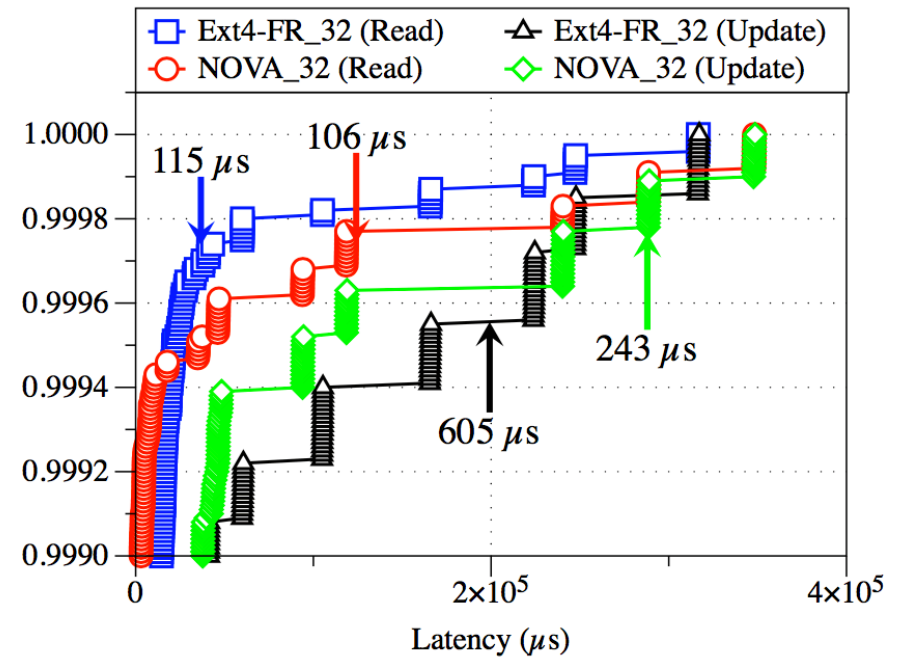
(b) YCSB-A

Performance evaluation

Comparison to NOVA-Fortis



(a) YCSB Workloads



(b) YCSB-A

* Issue with Linux implementation and performance reporting

- close() system call waits for background flush to complete
- even though, with F-Responder, no need to wait

F-Responder summary

- Reap PM performance through First Responder
- Despite using legacy file system and storage
- By background trekking of critical storage stack

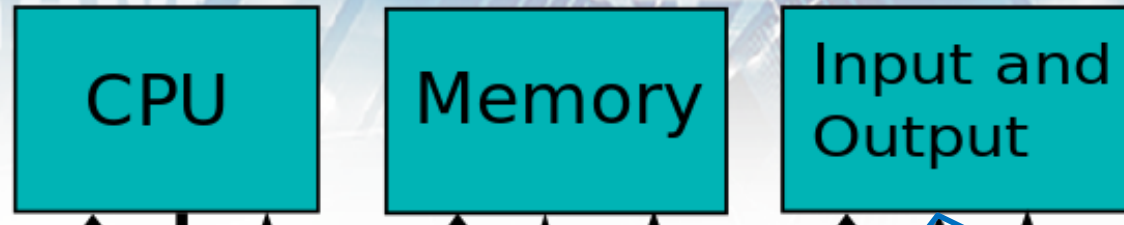
It's the storage stack, stupid!

Overall Summary

NV RAM



NOT! your father's ~~Oldsmobile~~ Storage Device



Types of Peripheral Devices



Peripheral: Auxiliary, Supplementary, relating to periphery

PAST storage topics of interest?

- **RAID**
 - Increase I/O bandwidth
- **Buffer Caching**
 - Improve latency
- **Swapping**
 - Improve resource sharing



**Revisit &
Rediscover**

Take a fresh look at these old favorites.

Thank you!!!



NECSST

Next-generation Embedded / Computer System Software Technology